Binary search trees of permuton samples

Benoît Corsini, Victor Dubach, Valentin Féray

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- I. BSTs and permutons
- II. Universality of the height
- III. Proof ideas
- IV. Subtree sizes

Binary Search Trees

A Binary Search Tree (BST) is a rooted labeled binary tree such that: for each vertex v, all labels of vertices in its left-subtree are smaller than that of v (resp. right-subtree, greater).

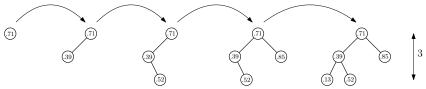
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BST associated with a sequence of distinct numbers

Let $y = (y_1, \ldots, y_n)$ be a sequence of distinct numbers. We can construct a BST $\mathcal{T}\langle y \rangle$ by successively adding leaves with those labels.

Construction of $\mathcal{T}\langle y \rangle$ with y = (.71, .39, .52, .85, .13):



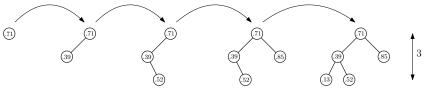
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Rem: same shape as $\mathcal{T}\langle\sigma\rangle$ with $\sigma = (4, 2, 3, 5, 1)$.

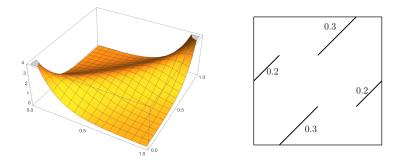
<u>Goal:</u> Study $\mathcal{T}\langle\sigma\rangle$ where σ is a random permutation.

Q: What kind of randomness?

Permutons

A *permuton* is a probability measure μ on $[0,1]^2$ with uniform marginals:

 $\forall t \in [0,1], \quad \mu([0,t] \times [0,1]) = \mu([0,1] \times [0,t]) = t.$



Permuton samples

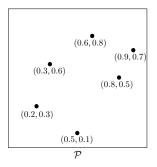
Let $\mathcal{P} = \{(x_1, y_1), \dots, (x_n, y_n)\} \subset \mathbb{R}^2$ with no common x- or y-coordinate.

Define a reordering $(x_{(1)}, y_{(1)}), \ldots, (x_{(n)}, y_{(n)})$ by $x_{(1)} < \cdots < x_{(n)}$. There exists a unique permutation $\sigma = \sigma \langle \mathcal{P} \rangle$ such that $(y_{(1)}, \ldots, y_{(n)})$ and $(\sigma(1), \ldots, \sigma(n))$ are in the same relative order.

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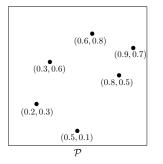


$$(y_{(1)}, \ldots, y_{(6)}) = (0.3, 0.6, 0.1, 0.8, 0.5, 0.7)$$

$$\sigma \langle \mathcal{P} \rangle = (2, 4, 1, 6, 3, 5)$$

Permuton samples

If \mathcal{P} is random i.i.d. under a permuton μ , this permutation σ_{μ}^{n} is a *permuton sample*.



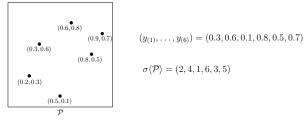
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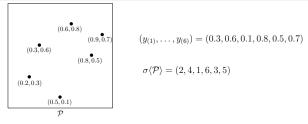


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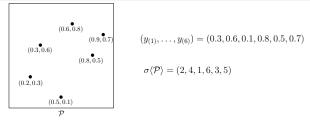


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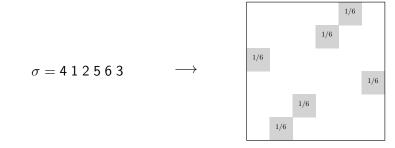
- Continuous deformation of the uniform distribution
- Highly non-parametric, wide range of distributions
- Insight on the links between random permutations and their limit permuton

More on permuton theory...

From permutations to permutons

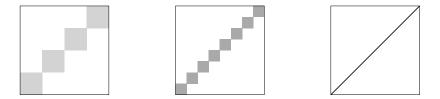
To each permutation $\sigma \in S_n$ we can associate a permuton μ_{σ} with:

$$\frac{d\mu_{\sigma}}{d\operatorname{Leb}_{[0,1]^2}} = n \cdot \sum_{i=1}^n \mathbb{1}_{\left[\frac{i-1}{n}, \frac{i}{n}\right] \times \left[\frac{\sigma(i)-1}{n}, \frac{\sigma(i)}{n}\right]}$$



Convergence of a sequence of permutations

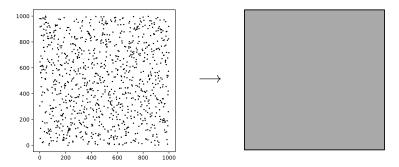
If $(\sigma_n)_{n\in\mathbb{N}}$ is a sequence of permutations and μ is a permuton, we say that $(\sigma_n)_{n\in\mathbb{N}}$ converges towards μ when $\mu_{\sigma_n} \xrightarrow[n\to\infty]{} \mu$ weakly.



Identity permutations converge to the diagonal permuton.

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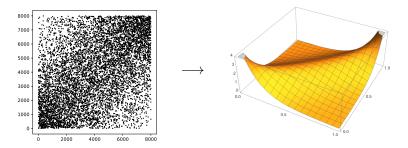
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Uniform permutations converge a.s. to Leb_{[0,1]²}.

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Mallows permutations with parameter $q_n = 1 - \frac{\beta}{n}$ (cf later) converge in probability to the Mallows permuton μ_{β} .

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Hoppen et al '13

For any permuton μ , the sequence of sampled permutations (σ_{μ}^{n}) converges a.s. to μ .

Goal (reminder):

Study $\mathcal{T}\langle \sigma_{\mu}^{n} \rangle$ where σ_{μ}^{n} is sampled from a permuton μ .

BST of a point process

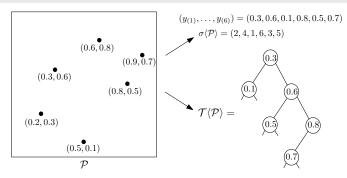
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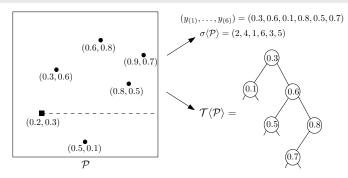


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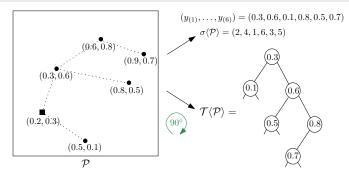


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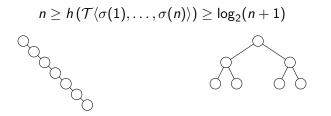
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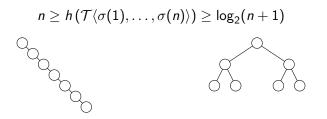
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Depending on the law of σ , asymptotic results for $h(\mathcal{T}\langle\sigma\rangle)$?

Devroye '86

If $\sigma^n \sim \text{Unif}(\mathfrak{S}_n)$ then, as $n \to \infty$: $\frac{h(\mathcal{T}\langle \sigma^n \rangle)}{c^* \log n} \longrightarrow 1$ in probability and in L^p for all $p \ge 1$, where $c^* \ge 2$ solves $c \log(2e/c) = 1$.

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The root is labeled by $\sigma(1) \sim \text{Unif}([n])$; the left-subtree has size $\sigma(1) - 1$ and the right-subtree has size $n - \sigma(1)$.

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At each node, the split between the left- and right-subtree is uniform \rightarrow the height is related to products of U_i or $1 - U_i$, where the U_i 's are i.i.d. uniform variables in [0, 1].

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Let $q \in [0, 1]$. The Mallows distribution $M_{n,q}$ is defined by: $\forall \sigma \in \mathfrak{S}_n, \quad M_{n,q}(\sigma) \propto q^{\mathrm{inv}(\sigma)}.$

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Addario-Berry - Corsini '21

If
$$\sigma^n \sim M_{n,q_n}$$
 where $n(1-q_n)/\log n \to 0$, then:

$$\frac{h(\mathcal{T}\langle \sigma^n \rangle)}{c^*\log n} \longrightarrow 1$$
in probability and in L^p for all $p > 1$.

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Let $\theta > 0$. The record-biased distribution $R_{n,\theta}$ is defined by:

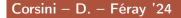
$$\forall \sigma \in \mathfrak{S}_n, \quad R_{n,\theta}(\sigma) \propto \theta^{\operatorname{rec}(\sigma)}.$$

Corsini '23

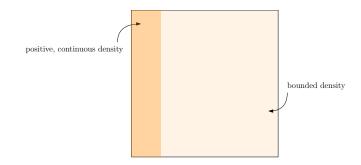
If $\sigma^n \sim R_{n,\theta}$ where $\theta \leq c^*$, then: $\frac{h(\mathcal{T}\langle \sigma^n \rangle)}{c^* \log n} \longrightarrow 1$ in probability and in L^p for all p > 1.

Assumption (A1):

 μ has a bounded density ρ on $[0, 1]^2$, which is continuous and positive on a neighborhood of $\{0\} \times [0, 1]$.

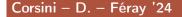


Let μ be a permuton satisfying (A1). Then: $\frac{h(\mathcal{T}\langle \sigma_{\mu}^{n} \rangle)}{c^{*} \log n} \longrightarrow 1$ in probability and in L^{p} for all $p \geq 1$.

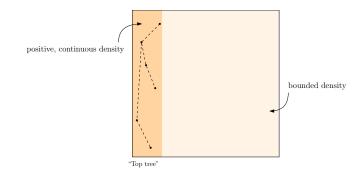


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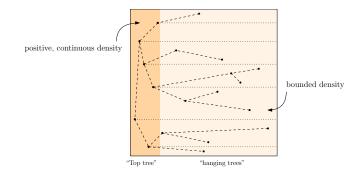


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Corsini – D. – Féray '24

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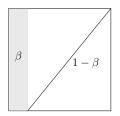
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Necessity of assumptions:



Height $\sim c_{\beta} \log n$ for some $c_{\beta} > c^*$



 $\text{Height} \ge c\sqrt{n}$

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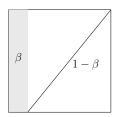
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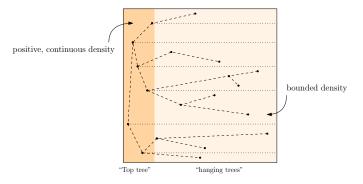
Question: Is there a permuton μ for which $\liminf \frac{h(\tau \langle \sigma_{\mu}^{n} \rangle)}{\log n} < c^{*}$?



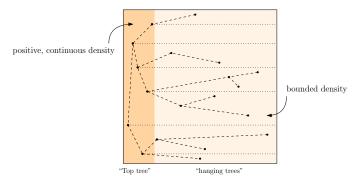
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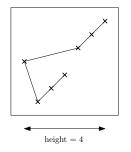
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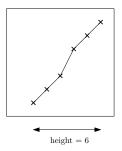
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Problem: taking out a point might double the height...





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Let $\mathcal{P}_{-} \subseteq \mathcal{P}_{+}$ be two point sets. If $C \subseteq \mathcal{P}_{+}$ is a chain in $\mathcal{T}\langle \mathcal{P}_{+} \rangle$, then $C \cap \mathcal{P}_{-}$ is a chain in $\mathcal{T}\langle \mathcal{P}_{-} \rangle$.

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$$h\left(\mathcal{T}\langle\mathcal{P}_{-}
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If $\mathcal{P}_{-} \subseteq \mathcal{P} \subseteq \mathcal{P}_{+}$, we can apply this lemma twice: with $\mathcal{P}_{-} \subseteq \mathcal{P}$, and with $\mathcal{P} \subseteq \mathcal{P}_{+}$. Thus:

$$h\left(\mathcal{T}\langle\mathcal{P}_+
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angle) - |\mathcal{C}_+ \cap (\mathcal{P}_+ \setminus \mathcal{P})| \le h\left(\mathcal{T}\langle\mathcal{P}
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negligible?

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1 "Poissonization": easier to work with a Poisson point process \mathcal{P} with intensity $n\rho$;

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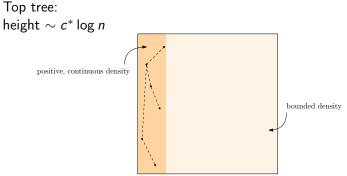
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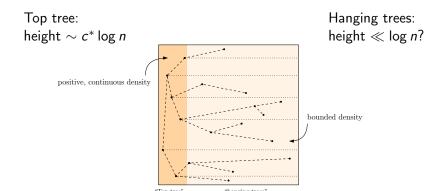


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Let $\{Y_i, 1 \le i \le n\}$ be i.i.d. random variables with law μ_0 on [0, 1], such that μ_0 has a positive density. Let $Y_{(1)} < \cdots < Y_{(n)}$ be their increasing reordering, and $\zeta_k := Y_{(k+1)} - Y_{(k)}$ be the *k*-th spacing. Then:

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 \longrightarrow comparison arguments along with deviation estimates.

Comparison with monotone subsequences

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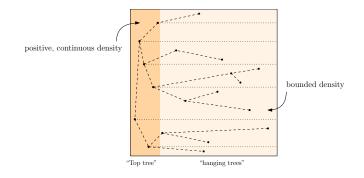
 \longrightarrow All hanging trees have height $\ll \log n$.

Assumption (A1):

 μ has a bounded density ρ on $[0, 1]^2$, which is continuous and positive on a neighborhood of $\{0\} \times [0, 1]$.

Corsini – D. – Féray '24

Let μ be a permuton satisfying (A1). Then: $\frac{h(\mathcal{T}\langle \sigma_{\mu}^{n} \rangle)}{c^{*} \log n} \longrightarrow 1$ in probability and in L^{p} for all $p \geq 1$.

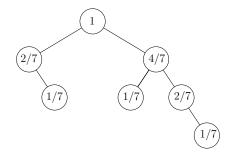


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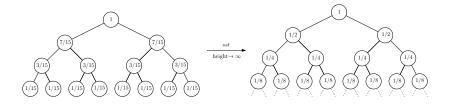
Subtree sizes

Let \mathcal{T} be a finite rooted tree. If $v \in \mathcal{T}$, we let $t(\mathcal{T}, v)$ be the proportion of nodes that are descendants of v (including itself).



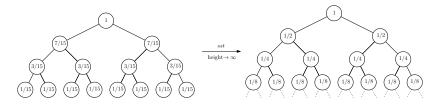
Subtree size convergence [Grübel '23]

Let (\mathcal{T}_n) be a sequence of binary trees. We say that it converges for the *subtree size topology* if $t(\mathcal{T}_n, v)$ converges for all fixed v. The decoration of the infinite binary tree with these limits is the *subtree size limit*.



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Rem: The limit object can be thought of as a probability measure on infinite paths.

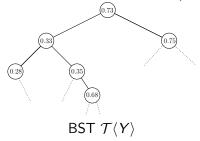
Subtree size limit sampled by a measure

Let *m* be a probability measure on [0,1] with no atom. Let $(Y_1, Y_2, ...)$ be i.i.d. under *m*, and let $\mathcal{T}\langle Y \rangle$ be their infinite BST. For each node *v*, let $\psi_m(v)$ be the difference between the labels of the leftmost ancestor to the right of *v*, and the rightmost ancestor to the left of *v*. Thus: ψ_m is a random subtree size limit.

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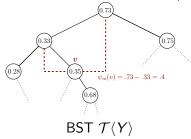
Ex: say m = 2x dx and Y = (0.73, 0.33, 0.75, 0.35, 0.68, 0.28, ...).



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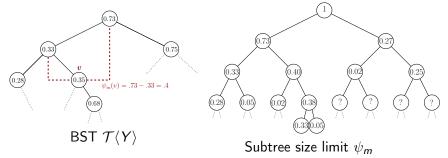
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Assumption (A2):

There exists a measure μ_0 on [0,1] with no atoms such that

$$\frac{1}{x}\mu([0,x]\times\cdot)\underset{x\to 0^+}{\longrightarrow}\mu_0.$$
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(A2) is weaker than (A1).

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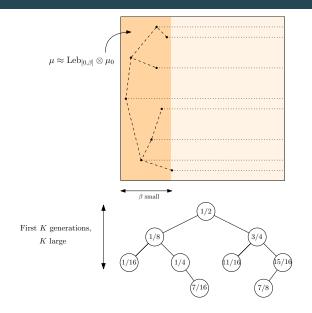
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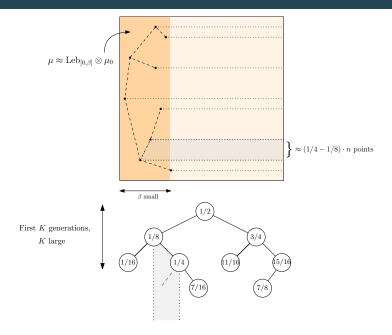
Corsini – D. – Féray '24

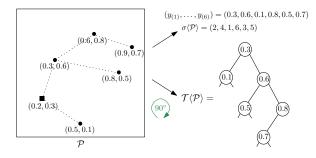
• If μ satisfies (A2) then, in distribution:

$$\mathcal{T}\langle \mathcal{P}^n_{\mu} \rangle \xrightarrow{\mathrm{ssc}} \psi_{\mu_0} .$$

If *T*(*P*ⁿ_μ) converges in distribution for the subtree size topology then (1) is satisfied, but μ₀ may have atoms.







Thank you for your attention!

