

Frobenius conjecture: Christoffel and Sturmian words

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Outline

We propose a generalization to the

Frobenius conjecture

using combinatorics on words.

We want to

- ▶ link the Frobenius conjecture with combinatorics of words
- ▶ propose a generalization of the conjecture
- ▶ prove it for the language of a Christoffel word and a Sturmian word

Frobenius conjecture

Markoff numbers

- ▶ **Markoff triple** : a solution of the Diophantine equation

$$x^2 + y^2 + z^2 = 3xyz$$

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$$1, 2, 5, 13, 29, 34, 89, 169, 194, \dots$$

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- ▶ $\mathcal{M}$: set of Markoff numbers

Markoff numbers

- Approximation theorem for irrational real numbers :

Theorem (Markoff, 1879-1880)

The *Lagrange spectrum* below 3 is given by

$$\mathcal{L}_{<3} = \left\{ \frac{\sqrt{9m^2 - 4}}{m} : m \in \mathcal{M} \right\}.$$

- Infimum of real indefinite binary quadratic form

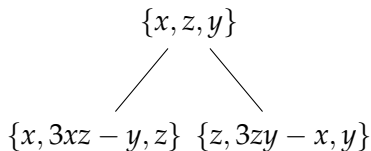
Theorem (Markoff, 1879-1880)

Let $f(x, y)$ be an indefinite binary quadratic form. Assume that its infimum $L(f)$ and its discriminant $d(f)$ satisfy the inequality $\sqrt{d(f)} < 3L(f)$. Then the infimum of f are attained and

$$\frac{\sqrt{d(f)}}{L(f)} = \sqrt{9 - 4\sqrt{m^2}}$$

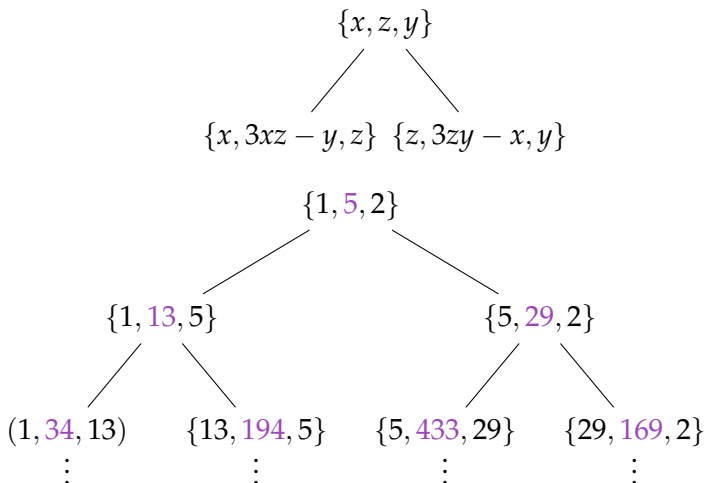
Markoff tree

- From a proper **Markoff triple**, we can construct two others distinct Markoff triples :



Markoff tree

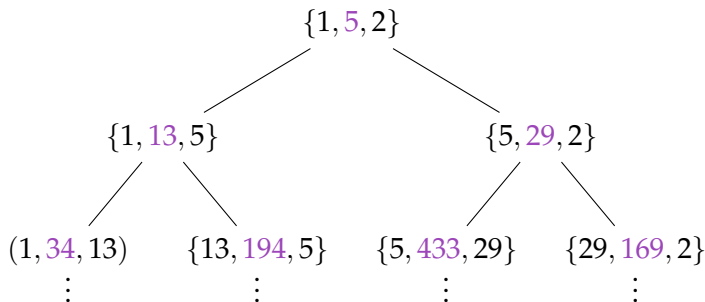
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Markoff tree

Theorem

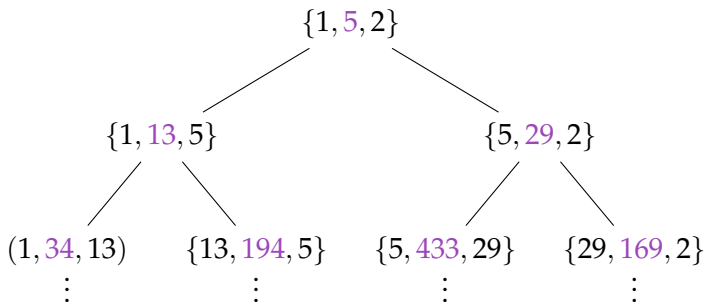
All proper Markoff triples appear exactly one in the Markoff tree.



Frobenius conjecture

Conjecture (Frobenius,1913)

*Every Markoff number appears exactly once as the **maximum** of a Markoff triple in the Markoff tree.*



Frobenius conjecture

- ▶ A lot of equivalent formulation of the conjecture are known (Aigner, 2013).
- ▶ Today, we will study this conjecture using combinatorics on words.

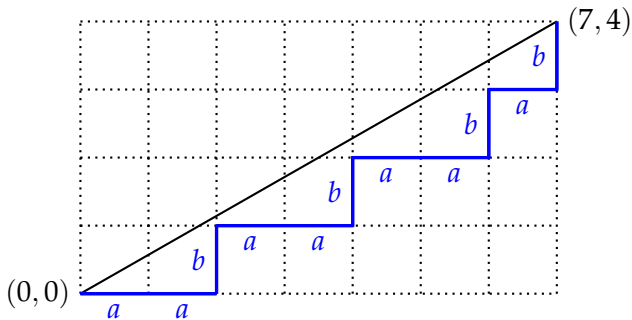
Christoffel words

Combinatorics on words

- ▶ **Word** : a sequence of element in a set $A = \{a, b\}$
- ▶ **Factor** : u is a factor of w if there exist two words x, y such that $w = xuy$.
- ▶ **Conjugate** : u and v are conjugates if there exist two words xy such that $u = xy$ and $v = yx$.
- ▶ **Circular factor** : u is a circular factor of w if it is a factor of a conjugate of w .
- ▶ **Palindrome** : A word equal to its reversal.

Christoffel words

Examples : The Christoffel words of slope $7/4$



Christoffel words in Markoff works

The continued fraction expansion of the Markoff constant

$$\sqrt{9 - \frac{4}{m^2}}, m \in \mathcal{M},$$

as an infinite words, have the form p^∞ , where p is a Christoffel word written on the alphabet $\{11, 22\}$.

Markoff numbers and Christoffel words

► $\mu : \{a, b\}^* \rightarrow \mathrm{SL}_2(\mathbb{N})$

$$\mu(a) = \begin{pmatrix} 2 & 1 \\ 1 & 1 \end{pmatrix} \quad \mu(b) = \begin{pmatrix} 5 & 2 \\ 2 & 1 \end{pmatrix}$$

► $M : \{a, b\}^* \rightarrow \mathbb{N} :$

$$M(w) = \mu(w)_{12}$$

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Theorem (Bombieri, 2007 ; Reutenauer, 2009)

The function f from the set of *Christoffel words* to the set of *Markoff triples* define as

$$f(w) = \{M(u), M(w), M(v)\},$$

where uv is the standard factorization of w , is a *bijection*.

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Conjecture (Frobenius)

The function M restricted to the set of Christoffel words is *injective*.

Other functions

- ▶ Thue-Morse morphism : $\theta(a) = ab, \theta(b) = ba$
- ▶ $w^{(+)} = xp\tilde{x}$, with p is the longest palindrome suffix of w
- ▶ Iterated palindromization : $Pal(\varepsilon) = \varepsilon$,

$$Pal(w_1 \dots w_n) = (Pal(w_1 \dots w_{n-1})w_n)^{(+)}$$

Theorem (Reutenauer and Vuillon, 2017)

For each word $v \in \{a, b\}^*$, the number

$$|Pal \circ \theta \circ Pal(av)| + 2$$

is a Markoff number. The mapping is surjective and the injection is equivalent to Frobenius conjecture.

Generalized Frobenius conjecture

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Conjecture

*The function M restricted to the set of **circular factors** of Christoffel words is injective.*

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Conjecture

*The function M restricted to the set of **circular factors** of Christoffel words is injective.*

- ▶ We don't have a proof of this conjecture.
- ▶ The generalized Frobenius conjecture implies the Frobenius conjecture.
- ▶ We can prove that the function M is **injective** on two subsets of the circular factors of Christoffel words.

Language of a Christoffel word

Proposition

*Let w be a Christoffel word and H_w the set of its circular factors. The function M restricted to H_w is *injective*.*

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Example : $H_{aaabaab}$

a	b	aa	ab	ba	aaa	aab	aba	baa	...
1	2	3	5	7	8	13	17	19	...

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Example : $H_{aaabaab}$

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Remark : In fact, we show that M is strictly increasing according to the radix order on H_w .

Language of a Christoffel word

w_1	a	a	a	b	a	a	b		1	3	8	34	115	311	1325
w_2	a	a	b	a	a	a	b		1	3	13	44	119	313	1327
w_3	a	a	b	a	a	b	a		1	3	13	44	119	507	1715
w_4	a	b	a	a	a	b	a	$\mapsto$	1	5	17	46	121	513	1735
w_5	a	b	a	a	b	a	a		1	5	17	46	196	663	1793
w_6	b	a	a	a	b	a	a		2	7	19	50	212	717	1939
w_7	b	a	a	b	a	a	a		2	7	19	81	274	741	1949

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1. Next line Lemma

Language of a Christoffel word

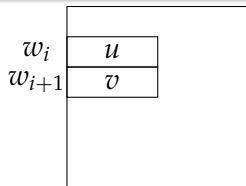
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1. Next line Lemma
2. Next column Lemma

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Lemma (Next line)

If u is a prefix of w_i and v is a prefix of w_{i+1} such that $|u| = |v|$ and $u \neq v$, then $M(u) < M(v)$.



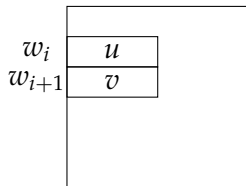
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Sketch of proof :

- $w_i = pabs$ and $w_{i+1} = pbas$
- Notice that sp is a palindrome

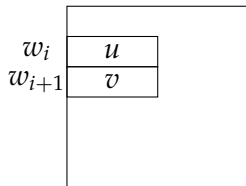


$a \ a \ \underline{a \ b} \ a \ a \ b$
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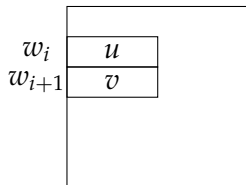
- ▶ $w_i = pabs$ and $w_{i+1} = pbas$
- ▶ Two cases :
 - ▶ $u = pa, v = pb$
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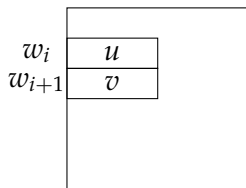
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 - ▶ $u = pab\tilde{p}t, v = pba\tilde{p}t$
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 - ▶ $u = pab\tilde{p}t, v = pb\tilde{a}p\tilde{t}$
 - ▶ $\mu(p) = \mu(\tilde{p})^t$
- ▶ Calculation of $M(v) - M(u) > 0$

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Lemma (Next column)

If u is a prefix of length k of w_n and v is a prefix of length $k + 1$ of w_1 , then $M(u) < M(v)$.

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If u is a prefix of length k of w_n and v is a prefix of length $k + 1$ of w_1 , then $M(u) < M(v)$.

Sketch of proof :

- ▶ Since w is a Christoffel word, we have that $w_1 = amb$ and $w_n = bma$, where m is a palindrome.
- ▶ $u = bp$ and $v = apx$, where p is a prefix of m and x is a letter.
- ▶ Let $\mu(p) = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$ be in the monoid $\mu(A^*)$.
- ▶ $M(v) - M(u) = (\mu(apx) - \mu(bp))_{12} = 2a - 3b + c - d > 0$

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Language of a Sturmian word

A **Sturmian word** is an infinite aperiodic word s such that $P(s, n) = n + 1$.

Corollary

Let S be a Sturmian words and $L(S)$ the set of all its factors. The function M restricted to $L(S)$ is injective.

Sketch of proof :

- ▶ Let u and v two factors of $L(S)$.
- ▶ There exists a prefix p of S such that u and v are both factors of p and p is a Christoffel word.

Remark

- Values of M are not increasing for the radix order in general

a	b	aa	ab	ba	bb	aaa	aab	aba	abb	baa
1	2	3	5	7	12	8	13	17	29	19

Reference

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