

Graph insertion operads

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Associative Algebra vs Operad

Associative Algebra A

extension cable



Operad $\mathcal{O}$

multi-socket



Associative Algebra vs Operad

Associative Algebra A

extension cable



product $\cdot$

words

Operad $\mathcal{O}$

multi-socket



partial composition $\circ_*$

trees

Associative Algebra vs Operad

Associative Algebra A

extension cable



product $\cdot$

words

$x \in A$

Operad $\mathcal{O}$

multi-socket



partial composition $\circ_*$

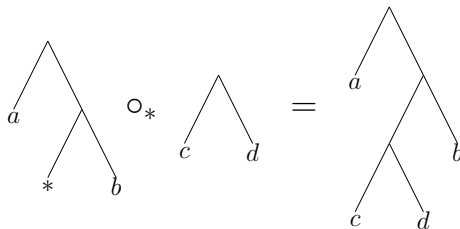
trees

$x \in \mathcal{O}[V]$

And now some examples

Example: Commutative magmatic

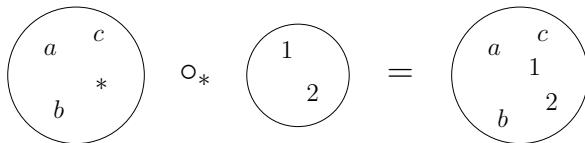
$\text{ComMag}[V]$ = non planar rooted binary trees with set of leaves V .



$$\text{ComMag}[\{a, b, *\}] \times \text{ComMag}[\{c, d\}] \rightarrow \text{ComMag}[\{a, b, c, d\}]$$

Example: Commutative

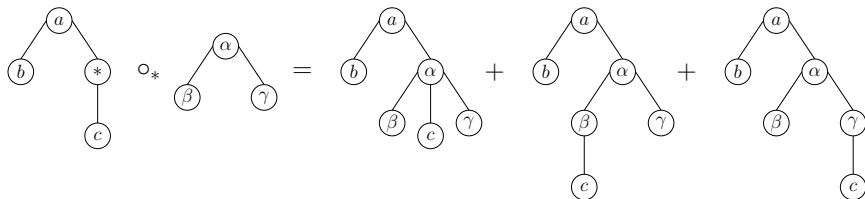
$$\mathbf{Com}[V] = \{V\}.$$



$$\mathbf{Com}[\{a, b, c, *\}] \times \mathbf{Com}[\{1, 2\}] \rightarrow \mathbf{Com}[\{a, b, c, 1, 2\}]$$

Example: Pre-Lie

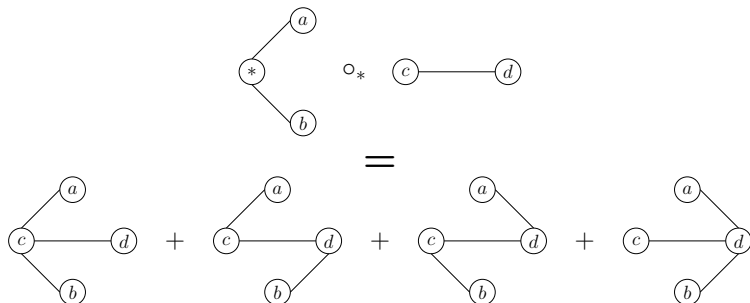
$\mathbf{PLie}[V]$ = non planar rooted trees with set of vertices V .



$$\mathbf{PLie}[\{a, b, c, *\}] \times \mathbf{PLie}[\{\alpha, \beta, \gamma\}] \rightarrow \mathbb{K}\mathbf{PLie}[\{a, b, c, \alpha, \beta, \gamma\}]$$

Example: Kontsevich-Willwacher

$\mathbf{G}[V]$ = simple graphs with set of vertices V .



$$\mathbf{G}[\{a, b, *\}] \times \mathbf{G}[\{c, d\}] \rightarrow \mathbb{K}\mathbf{G}[\{a, b, c, d\}]$$

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(Joyal) A *species* P is given by the data of:

for each finite set I , a vector space $\mathbb{K}P[I]$ of base $P[I]$,

for each bijection $\sigma : I \rightarrow J$, a linear map

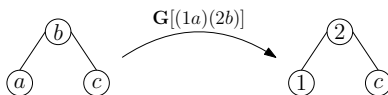
$$P[\sigma] : \mathbb{K}P[I] \rightarrow \mathbb{K}P[J],$$

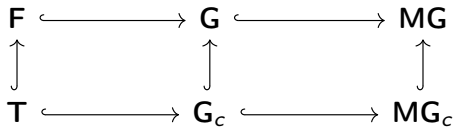
such that

$$P[\tau \circ \sigma] = P[\tau] \circ P[\sigma] \text{ and } P[\text{id}] = \text{id}.$$

$$\mathbf{G}[\{a, b, c\}] = \left\{ \begin{array}{c} \textcircled{b} \\ \textcircled{a} \quad \textcircled{c} \end{array}, \begin{array}{c} \textcircled{b} \\ \textcircled{a} \quad \textcircled{c} \end{array}, \begin{array}{c} \textcircled{b} \\ \textcircled{a} \text{---} \textcircled{c} \end{array}, \dots \right\}$$

$$3 \begin{array}{c} \textcircled{b} \\ \textcircled{a} \quad \textcircled{c} \end{array} - \begin{array}{c} \textcircled{b} \\ \textcircled{a} \quad \textcircled{c} \end{array} \in \mathbb{K}\mathbf{G}[\{a, b, c\}]$$





$$\begin{array}{ccccc}
 \mathbf{F} & \hookrightarrow & \mathbf{G} & \hookrightarrow & \mathbf{MG} \\
 \uparrow & & \uparrow & & \uparrow \\
 \mathbf{T} & \hookrightarrow & \mathbf{G}_c & \hookrightarrow & \mathbf{MG}_c
 \end{array}$$

pointing: $P^\bullet[V] = P[V] \times V$

$$\mathbf{T}^\bullet \hookrightarrow \mathbf{G}_c^\bullet \hookrightarrow \mathbf{MG}_c^\bullet$$

$$\begin{array}{ccccc}
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pointing: $P^\bullet[V] = P[V] \times V$

$$\mathbf{T}^\bullet \hookrightarrow \mathbf{G}_c^\bullet \hookrightarrow \mathbf{MG}_c^\bullet$$

$p \in P$ instead of “ V finite set and $p \in P[V]$ ”
 $x \in p$ if $p \in P[\{x, \dots\}]$

Partial composition

Let P be a species. A *partial composition* over P is given by:

for each disjoint pair of finite sets V_1, V_2

and for each element $* \in V_1$,

a map

$$\circ_* : P[V_1] \times P[V_2] \rightarrow \mathbb{K}P[V_1 \sqcup V_2 - \{*\}].$$

Partial composition

Let P be a species. A *partial composition* over P is given by:

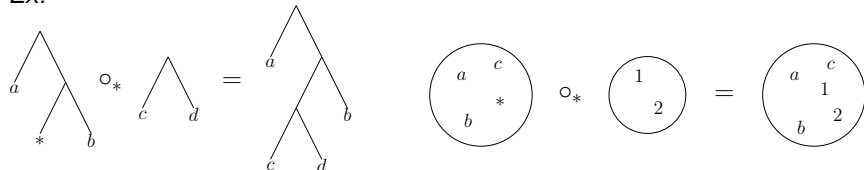
for each disjoint pair of finite sets V_1, V_2

and for each element $* \in V_1$,

a map

$$\circ_* : P[V_1] \times P[V_2] \rightarrow \mathbb{K}P[V_1 \sqcup V_2 - \{*\}].$$

Ex:



A partial composition is *associative* if, for every $p, q, r \in P$:

for $*_1 \in p$ and $*_2 \in q$:

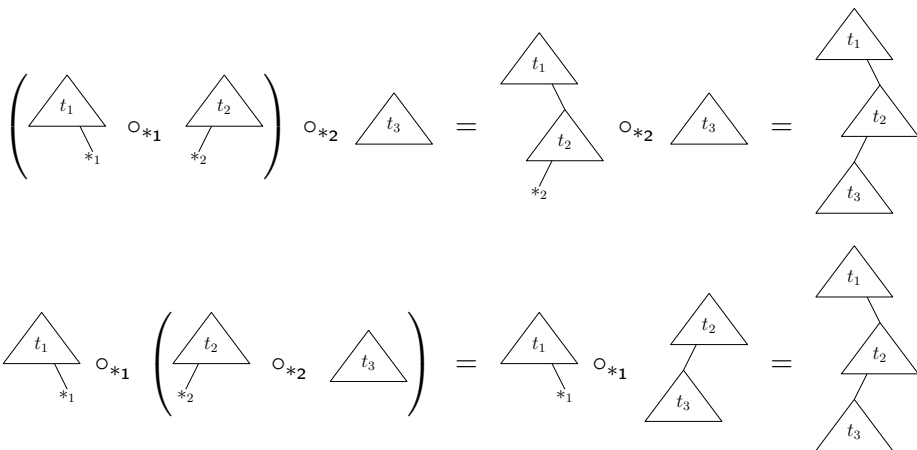
$$(p \circ_{*_1} q) \circ_{*_2} r = p \circ_{*_1} (q \circ_{*_2} r)$$

and for $*_1, *_2 \in p$:

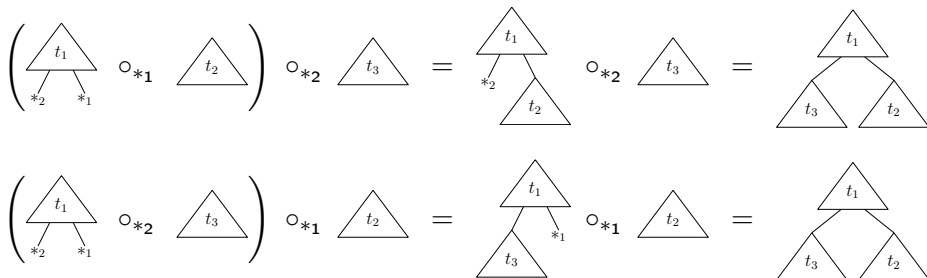
$$(p \circ_{*_1} q) \circ_{*_2} r = (p \circ_{*_2} r) \circ_{*_1} q$$

Partial composition

Ex: ComMag



Ex: ComMag



Definition

An *operad* is a species $\mathcal{O}$ with an associative partial composition $\circ$.

Ex: **Com**, **ComMag**, **NAP**, **PLie**, **Dend**, **Dup**, **TSig_k^{||}**...

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What about graphs ?

In physics (Connes, Kreimer), other branch of mathematics (Kontsevich-Willwacher).

Definition

An *operad* is a species $\mathcal{O}$ with an associative partial composition $\circ$.

Ex: **Com**, **ComMag**, **NAP**, **PLie**, **Dend**, **Dup**, **TSig_k^{||}**...

What about graphs ?

In physics (Connes, Kreimer), other branch of mathematics (Kontsevich-Willwacher).

$$\mathbf{PLie} \xrightarrow{\quad ? \quad} \mathbf{G}$$

Definition

An *operad* is a species $\mathcal{O}$ with an associative partial composition $\circ$.

Ex: **Com**, **ComMag**, **NAP**, **PLie**, **Dend**, **Dup**, **TSig_k^{||}**...

What about graphs ?

In physics (Connes, Kreimer), other branch of mathematics (Kontsevich-Willwacher).

PLie $\xrightarrow{\quad ? \quad}$ **MG**

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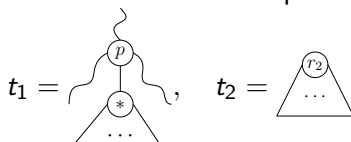
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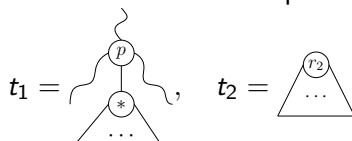
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Pre-Lie operad

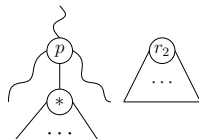
As a species: $\mathbf{PLie} = \mathbf{T}^\bullet$. Partial composition, $t_1 \circ_* t_2 = ?$



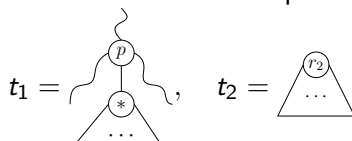
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1. take the disjoint union of t_1 and t_2

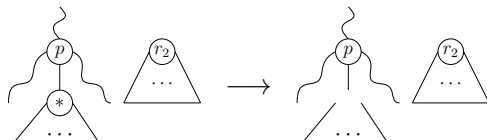


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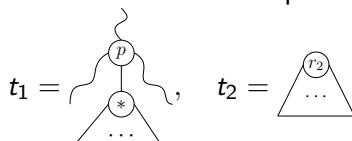


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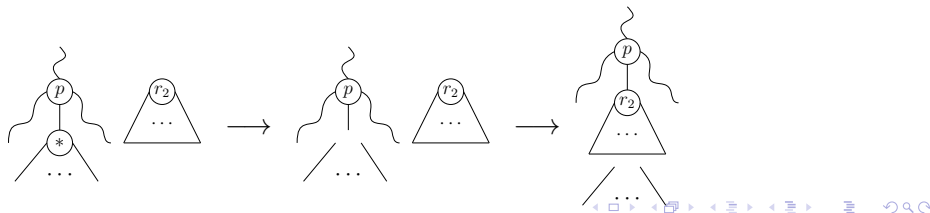
2. remove the vertex $*$



As a species: $\mathbf{PLie} = \mathbf{T}^\bullet$. Partial composition, $t_1 \circ_* t_2 = ?$

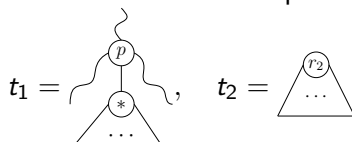


1. take the disjoint union of t_1 and t_2
2. remove the vertex $*$
3. if $* \neq r_1$, connect the parent of $*$ to r_2

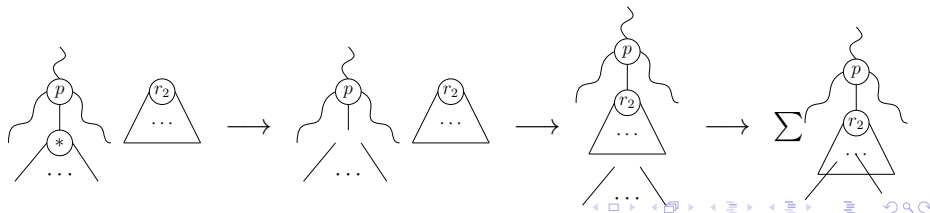


Pre-Lie operad

As a species: $\mathbf{PLie} = \mathbf{T}^\bullet$. Partial composition, $t_1 \circ_* t_2 = ?$

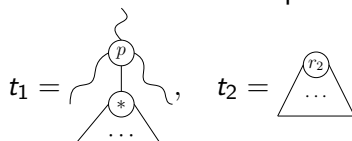


1. take the disjoint union of t_1 and t_2
2. remove the vertex $*$
3. if $* \neq r_1$, connect the parent of $*$ to r_2
4. sum over all the ways to connect the children of $*$ to t_2

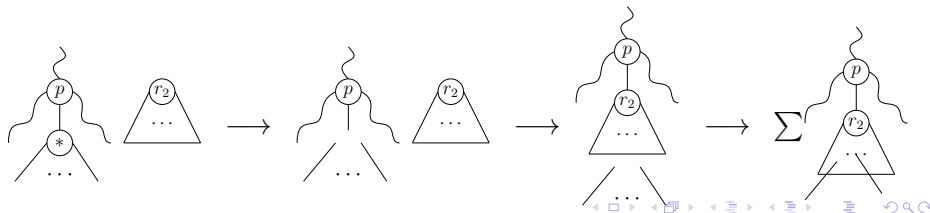


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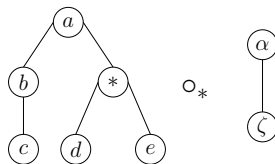


1. take the disjoint union of t_1 and t_2
2. remove the vertex $*$
3. if $* \neq r_1$, connect the parent of $*$ to r_2
4. sum over all the ways to connect the children of $*$ to t_2
5. the new root is r_1 if $* \neq r_1$ and r_2 else.

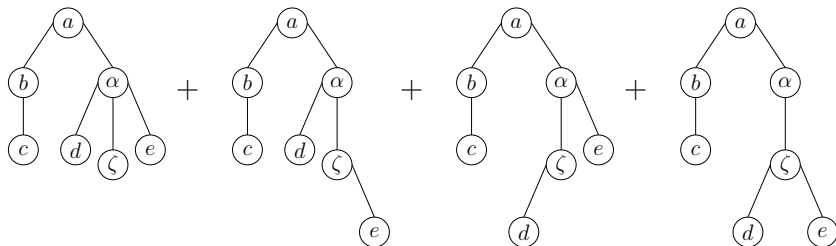


Pre-Lie operad

Ex:



=



As a species: **MG**. Partial composition, $g_1 \circ_* g_2 = ?$

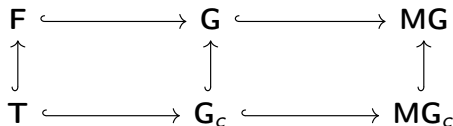
1. take the disjoint union of g_1 and g_2
2. remove the vertex $*$
3. sum over all the ways to connect the neighbours of $*$ to g_2 .

Canonical multigraph operad

As a species: **MG**. Partial composition, $g_1 \circ_* g_2 = ?$

1. take the disjoint union of g_1 and g_2
2. remove the vertex $*$
3. sum over all the ways to connect the neighbours of $*$ to g_2 .

Remark:



On **G**: Kontsevich-Willwacher.

Canonical multigraph operad

Ex:

$$\begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{*} \end{array} \circ_* \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} =$$

$$\begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} + \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} + 2 \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} \\
 + 2 \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} + 2 \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} + 4 \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} \\
 + \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} + \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array} + 2 \begin{array}{c} \text{a} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{b} \end{array} \begin{array}{c} \text{---} \end{array} \begin{array}{c} \text{c} \end{array}.$$

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As species: $\mathbf{T}^\bullet \hookrightarrow \mathbf{MG}_c^\bullet$

As operads: $\mathbf{PLie} \hookrightarrow \mathbf{MG}_c^\bullet ?$

The ϕ morphism

$$\begin{aligned}\phi : \mathbf{MG}_c &\hookrightarrow \mathbf{MG}_c^\bullet \\ g &\mapsto \sum_{r \in g} (g, r),\end{aligned}$$

Proposition

The species morphism ϕ induces an operad morphism $\mathbf{T} \hookrightarrow \mathbf{PLie}$.

The ϕ morphism

$$\begin{aligned}\phi : \mathbf{MG}_c &\hookrightarrow \mathbf{MG}_c^\bullet \\ g &\mapsto \sum_{r \in g} (g, r),\end{aligned}$$

Proposition

The species morphism ϕ induces an operad morphism $\mathbf{T} \hookrightarrow \mathbf{PLie}$.

$$\begin{array}{ccc}\mathbf{T} & \xrightarrow{\phi} & \mathbf{PLie} \\ \downarrow & & \downarrow \\ \mathbf{MG}_c & \xrightarrow{\phi} & \mathbf{MG}_c^\bullet\end{array}$$

The ϕ morphism

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Proposition

The species morphism ϕ induces an operad morphism $\mathbf{T} \hookrightarrow \mathbf{PLie}$.

$$\begin{array}{ccccc}\mathbf{T} & \xrightarrow{\sim} & \phi(\mathbf{T}) & \hookrightarrow & \mathbf{PLie} \\ \downarrow & & \downarrow & & \downarrow \\ \mathbf{MG}_c & \xrightarrow{\sim} & \phi(\mathbf{MG}_c) & \hookrightarrow & \mathbf{MG}_c^\bullet\end{array}$$

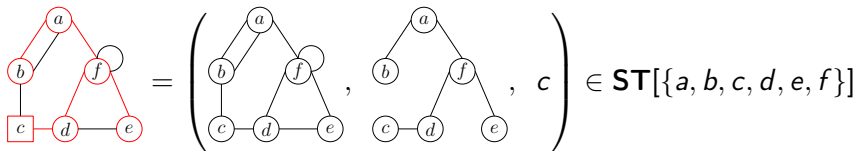
Spanning tree species

~~$\mathbf{MG}_c^\bullet$~~

→ **ST**: spanning tree species.

$$\mathbf{ST}[V] = \{(g, t, v) \mid g \in \mathbf{MG}_c[V], t \text{ spanning tree of } g, v \in V\}$$

Ex:



Definition (Hadamard product)

The Hadamard product $P \times Q$ of two species is defined by $P \times Q[V] = P[V] \times Q[V]$.

If P and Q are operads then $P \times Q$ is also an operad with $(p_1, q_1) \circ_* (p_2, q_2) = (p_1 \circ_* p_2, q_1 \circ_* q_2)$.

Definition (Hadamard product)

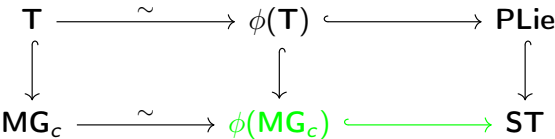
The Hadamard product $P \times Q$ of two species is defined by $P \times Q[V] = P[V] \times Q[V]$.

If P and Q are operads then $P \times Q$ is also an operad with $(p_1, q_1) \circ_* (p_2, q_2) = (p_1 \circ_* p_2, q_1 \circ_* q_2)$.

$$\mathbf{ST} \xrightarrow{\sim} \mathbf{MG} \times \mathbf{PLie}$$

$$(g, t, v) \mapsto (g \setminus t, (t, v))$$

$\Rightarrow$ operad structure on $\mathbf{ST}$



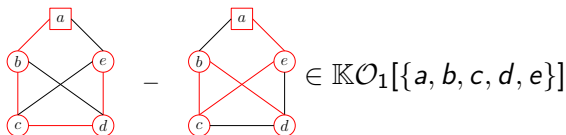
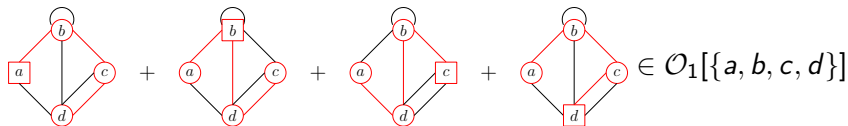
The sub-operad $\mathcal{O}_1$

$$\mathcal{O}_1[V] = \left\{ \sum_{v \in V} (g, t_v, v) \mid g \in MG_c[V], t_v \text{ spanning tree of } g \right\} \subset \mathbb{K}\mathbf{ST}[V]$$

The sub-operad $\mathcal{O}_1$

$$\mathcal{O}_1[V] = \left\{ \sum_{v \in V} (g, t_v, v) \mid g \in MG_c[V], t_v \text{ spanning tree of } g \right\} \subset \mathbb{KST}[V]$$

Ex:



The sub-operad $\mathcal{O}_1$

Lemma

$\mathcal{O}_1$ is a sub-operad of **ST** (i.e $\mathcal{O}_1 \circ_* \mathcal{O}_1 \subset \mathcal{O}_1$).

$$\mathcal{O}_1 \cap \mathbf{PLie} = \phi(T)$$

The sub-operad $\mathcal{O}_1$

Lemma

$\mathcal{O}_1$ is a sub-operad of **ST** (i.e $\mathcal{O}_1 \circ_* \mathcal{O}_1 \subset \mathcal{O}_1$).

$$\mathcal{O}_1 \cap \mathbf{PLie} = \phi(T)$$

$$\text{but } \mathcal{O}_1 \neq \phi(\mathbf{MG}).$$

$$\phi(\mathbf{MG}) \quad \mathcal{O}_1$$

$$\sum_{v \in g} (g, v) \quad \sum_{v \in g} (g, t_v, v)$$

How to choose a spanning tree for each vertex ?

$$\phi(\mathbf{MG}) \quad \mathcal{O}_1$$

$$\sum_{v \in g} (g, v) \quad \sum_{v \in g} (g, t_v, v)$$

How to choose a spanning tree for each vertex ?

We don't need to choose!

Definition

An ideal of an operad $\mathcal{O}$ is a sub-species $\mathcal{I}$ such that $\mathcal{O} \circ_* \mathcal{I} \subseteq \mathcal{I}$ and $\mathcal{I} \circ_* \mathcal{O} \subseteq \mathcal{I}$.

The quotient species $\mathcal{O}/\mathcal{I}$ defined by $\mathbb{K}\mathcal{O}/\mathcal{I}[V] = \mathbb{K}\mathcal{O}[V]/\mathbb{K}\mathcal{I}[V]$ then has a natural operad structure.

Ex:

$$\mathcal{O} = \mathbf{ComMag} \quad \mathcal{I} = \left\langle \begin{array}{c} \diagup \quad \diagdown \\ \diagdown \quad \diagup \\ a \quad b \end{array} \quad c - \begin{array}{c} \diagup \quad \diagdown \\ \diagup \quad \diagdown \\ b \quad c \end{array} \quad a \right\rangle \quad \mathbf{ComMag}/\mathcal{I} \cong \mathbf{Com}$$

The ideal $\mathcal{I}_1$

$$\begin{aligned}\mathcal{I}_1[V] = \{ & (g, t_1, v) - (g, t_2, v) \mid g \in \mathbf{MG}_c[V], \\ & t_1, t_2 \text{ spanning trees of } g, v \in V\} \\ & \subset \mathbb{K}\mathbf{ST}[V]\end{aligned}$$

Lemma

$\mathcal{I}_1$ is an ideal of $\mathcal{O}_1$.

$$\mathcal{O}_1/\mathcal{I}_1 \cong \phi(\mathbf{MG}_c)$$

Complete diagram

$$\begin{array}{ccccccc}
 \mathbf{T} & \xrightarrow{\sim} & \mathbf{PLie} \cap \mathcal{O}_1 / \mathcal{I}_1 & \equiv & \mathbf{PLie} \cap \mathcal{O}_1 & \hookrightarrow & \mathbf{PLie} \\
 \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 \mathbf{MG}_c & \xrightarrow{\sim} & \mathcal{O}_1 / \mathcal{I}_1 & \ll & \mathcal{O}_1 & \hookrightarrow & \mathbf{ST}
 \end{array}$$

Complete diagram

$$\begin{array}{ccccccc}
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 \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 \mathbf{G}_c & \xrightarrow{\sim} & \mathbf{G}_c \cap \mathcal{O}_1 / \mathcal{I}_1 & \ll & \mathbf{G}_c \cap \mathcal{O}_1 & \hookrightarrow & \mathbf{G}_c \cap \mathbf{ST} \\
 \downarrow & & \downarrow & & \downarrow & & \downarrow \\
 \mathbf{MG}_c & \xrightarrow{\sim} & \mathcal{O}_1 / \mathcal{I}_1 & \ll & \mathcal{O}_1 & \hookrightarrow & \mathbf{ST}
 \end{array}$$

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- Intuition and examples
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2 Extending Pre-Lie

- Pre-Lie and multigraph operads
- Link and extension

3 Finitely generated sub-operads

Generators and relations

Description of an operad by its generators and relations.

Ex: **PLie**

generators: $\left\{ \begin{array}{c} \boxed{a} \\ | \\ \bigcirc b \end{array} \right\}.$

relations:

$$\begin{array}{c} \boxed{*} \\ | \\ \bigcirc c \end{array} \circ_* \begin{array}{c} \boxed{a} \\ | \\ \bigcirc b \end{array} - \begin{array}{c} \boxed{a} \\ | \\ \bigcirc * \end{array} \circ_* \begin{array}{c} \boxed{b} \\ | \\ \bigcirc c \end{array} = \begin{array}{c} \boxed{*} \\ | \\ \bigcirc b \end{array} \circ_* \begin{array}{c} \boxed{a} \\ | \\ \bigcirc c \end{array} - \begin{array}{c} \boxed{a} \\ | \\ \bigcirc * \end{array} \circ_* \begin{array}{c} \boxed{c} \\ | \\ \bigcirc b \end{array}$$

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What about **ST** ? $\mathbf{ST} \cong \mathbf{MG} \times \mathbf{PLie}$

→ what about **MG** ?

Generators and relations

MG: too involved.

MG_c ?

Generators and relations

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G_c ?

Generators and relations

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... **T** ?

Generators and relations

MG: too involved.

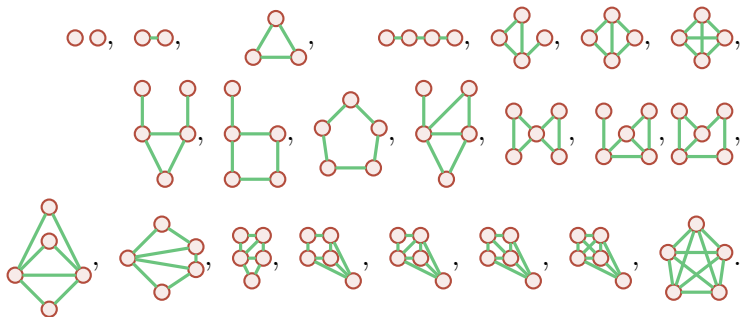
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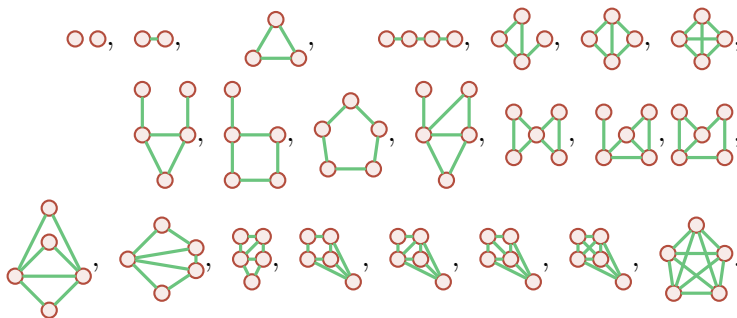
... **T** ? still too involved.

→ finitely generated sub-operads.

Graph generators



Graph generators



Lemma

Any graph with n vertices and at least $\binom{n-1}{2} + 1$ edges must be a generator.

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Proof: $v(g)$ = number of vertices, $e(g)$ number of edges.

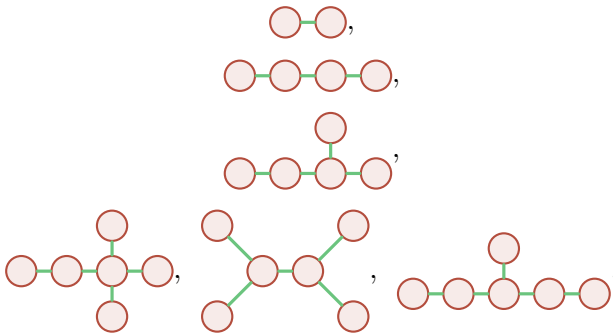
$$g_1 \circ_* g_2 = g + \dots \Rightarrow e(g) = e(g_1) + e(g_2)$$

$$\text{maximum when } e(g_i) = \binom{v(g_i)}{2}$$

$$\binom{v(g_1)}{2} + \binom{v(g_2)}{2} = \binom{v(g_1)}{2} + \binom{v(g)+1-v(g_1)}{2}$$

$$\text{maximum when } v(g) = 2 \text{ or } v(g) = n - 1 \text{ and equal to } \binom{n-1}{2}.$$

Tree generators



finitely generated sub-operads

- $\{ \textcircled{a} \textcircled{b} \}$
relation:

$$\textcircled{a} \textcircled{*} \circ_* \textcircled{b} \textcircled{c} = \textcircled{*} \textcircled{c} \circ_* \textcircled{a} \textcircled{b}.$$

finitely generated sub-operads

- $\{\circledast a \circledast b\} \cong \mathbf{Com}$
relation:

$$\circledast a \circledast * \circledast b \circledast c = \circledast * \circledast c \circledast * \circledast a \circledast b.$$

finitely generated sub-operads

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- $\{\textcircled{a} \text{---} \textcircled{b}\} \cong \mathbf{ComMag}$

no relations.

- SP

generators: $\{ \textcircled{a} \textcircled{b}, \textcircled{a} \text{---} \textcircled{b} \}$

relations:

$$\textcircled{a} \textcircled{*} \textcircled{b} \textcircled{c} = \textcircled{*} \textcircled{c} \textcircled{a} \textcircled{b},$$

$$\textcircled{a} \text{---} \textcircled{*} \textcircled{b} \textcircled{c} = \textcircled{c} \textcircled{*} \textcircled{a} \text{---} \textcircled{b} + \textcircled{b} \textcircled{*} \textcircled{a} \text{---} \textcircled{c}.$$

- **SP Koszul.**

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- **SP[!]** generators: $\{ \textcircled{a} \textcircled{b}^{\vee}, \textcircled{a} \text{---} \textcircled{b}^{\vee} \}$

relations:

$$\textcircled{a} \text{---} \textcircled{*}^{\vee} \circ_* \textcircled{b} \textcircled{c}^{\vee} = 0,$$

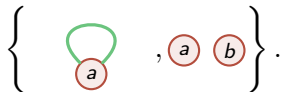
$$\textcircled{a} \textcircled{*}^{\vee} \circ_* \textcircled{b} \text{---} \textcircled{c}^{\vee} + \textcircled{c} \text{---} \textcircled{*}^{\vee} \circ_* \textcircled{a} \textcircled{b}^{\vee} + \textcircled{b} \text{---} \textcircled{*}^{\vee} \circ_* \textcircled{a} \textcircled{c}^{\vee} = 0,$$

$$\textcircled{a} \textcircled{*}^{\vee} \circ_* \textcircled{b} \textcircled{c}^{\vee} + \textcircled{c} \textcircled{*}^{\vee} \circ_* \textcircled{a} \textcircled{b}^{\vee} + \textcircled{b} \textcircled{*}^{\vee} \circ_* \textcircled{c} \textcircled{a}^{\vee} = 0.$$

finitely generated sub-operads

- LP

generators:



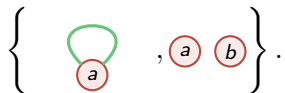
relations:

$$\begin{array}{c} (a) * (b) \circ_* (c) = (*)(c) \circ_* (a)(b) \\ \dots ? \end{array}$$

finitely generated sub-operads

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relations:

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We have $\mathbf{SP} \subset \mathbf{LP}$ but $a \text{---} b \text{---} c \notin \mathbf{LP}$.

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$a \text{---} b \text{---} c \notin \mathbf{LP}$: Computer algebra

Extensions and links of **PLie** and Kontsevich-Willwacher operads.
Study of finite case.

Conclusion

Extensions and links of **PLie** and Kontsevich-Willwacher operads.
Study of finite case.
General method to compute graph operads.

Extensions and links of **PLie** and Kontsevich-Willwacher operads.

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Perspectives

generators and relations of **ST** and co ?

description of **LP** ?

SP[!] link with permutations ?

ComMag $\hookrightarrow$ **T** $\hookrightarrow$ **PLie** further factorization ?

Thank you for your attention.

arXiv:1912.06563