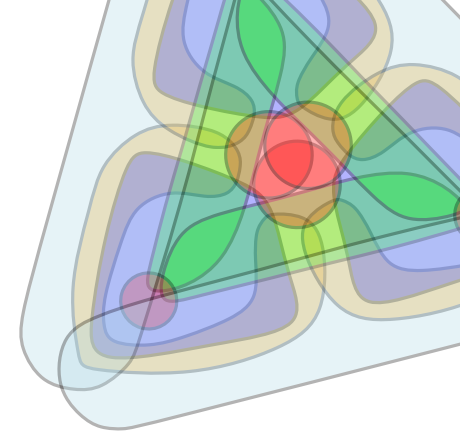
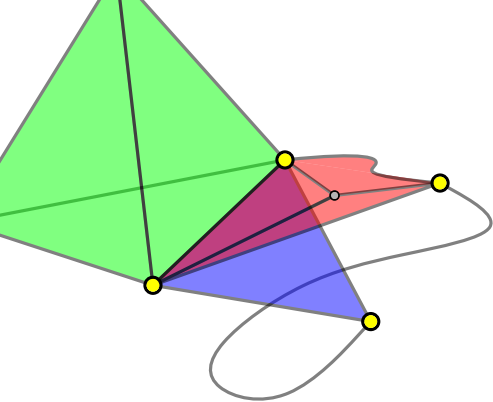
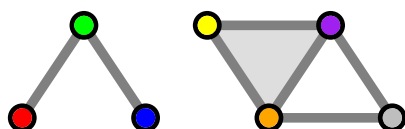
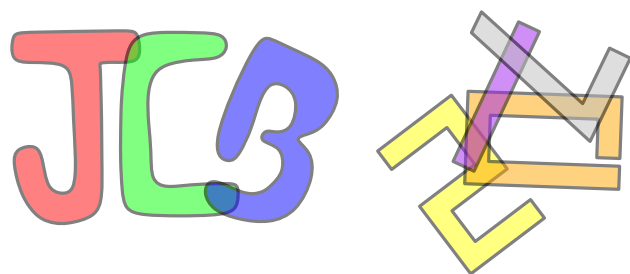




Helly-type theorem and topological combinatorics



Xavier Goaoc

"topological combinatorics" ...

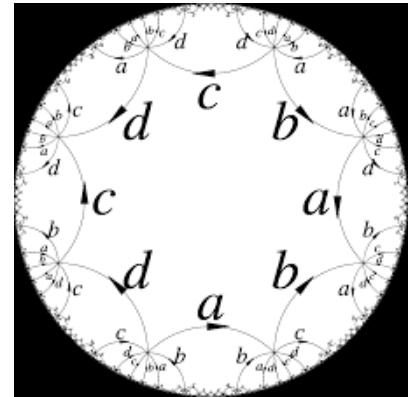
... what do **combinatorics** and **topology** have to do with one another?

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Combinatorics routinely used in topology.

Discrete approaches (eg. Whitehead's collapses)

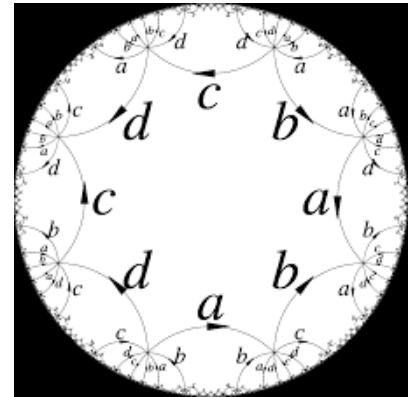


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The use of topology in combinatorics is perhaps more surprising.

Győry and Lovász (1977-1978): any k -connected graph $G = (V, E)$ can be partitioned into k subsets that induce connected subgraph, have prescribed size (summing to $|V|$) and each contain a prescribed element.

Lovász (1978): if the n -elements subsets of the $\{1, 2, \dots, 2n + k\}$ are partitioned into $k + 1$ classes then some class contains disjoint subsets.

Principle of a proof of the Kneser conjecture.

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Chromatic number of the graph $G = (V, E)$ where $V = \binom{[2n+k]}{n}$ and E connects disjoint subsets.

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Equivalent: For any continuous map $f : \mathbb{S}^d \rightarrow \mathbb{R}^d$ there exists $x \in \mathbb{S}^d$ such that $f(x) = f(-x)$.

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Blueprint:

Associate to G some topological space X with an antipodality action via some "neighborhood" simplicial complex.

Translate c -colorings of G into antipodal maps $X \rightarrow \mathbb{S}^{c-2}$.

Bound the connectivity of X and apply the Borsuk-Ulam theorem.

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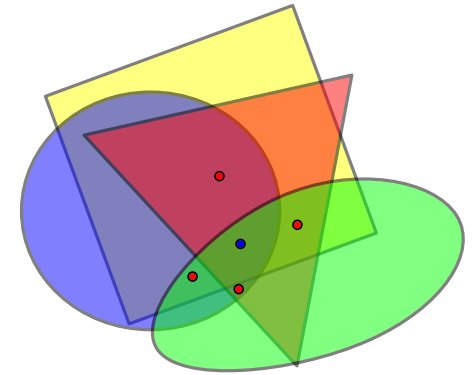
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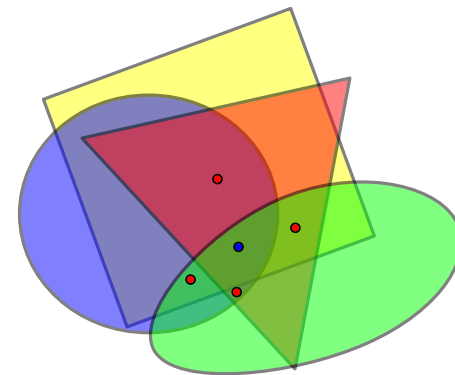
Examine similar approaches, for geometric questions, in a few more words...

Helly's Theorem. Any finite family of convex sets in $\mathbb{R}^d$ has non-empty intersection if every $d + 1$ elements have non-empty intersection.



[Helly, 1913]

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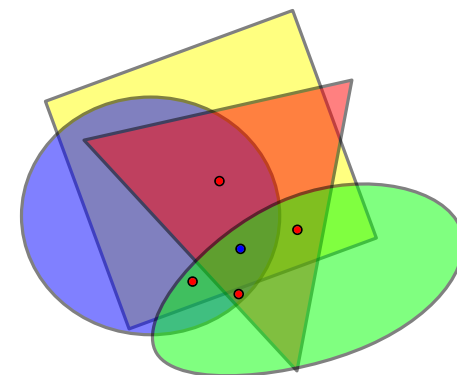
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This talk:

Examine Helly's theorem via **simplicial complexes**.

Relax "convexity" into more general topological conditions.

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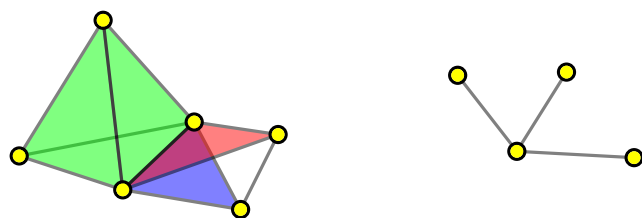
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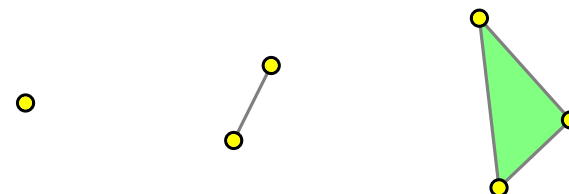
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geometric simplicial complex

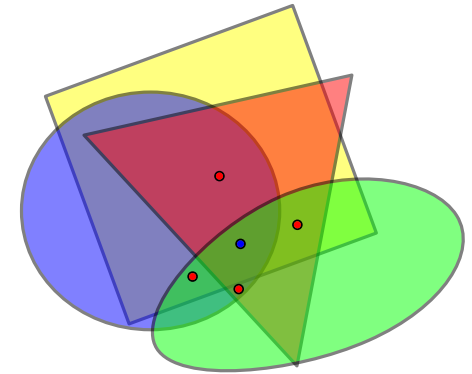
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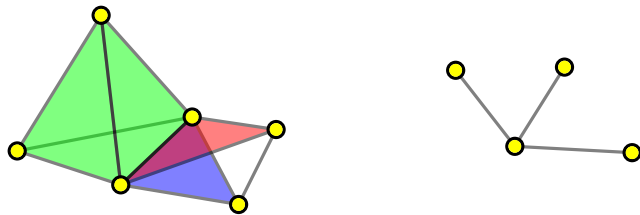
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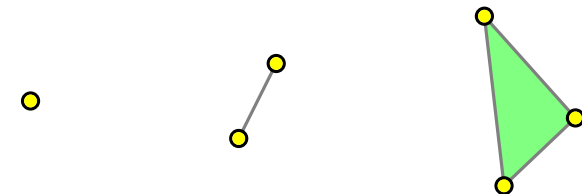


abstract simplicial complex

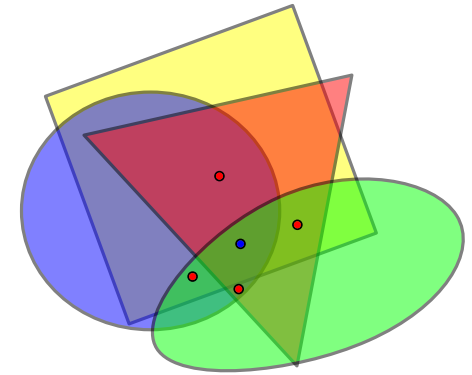
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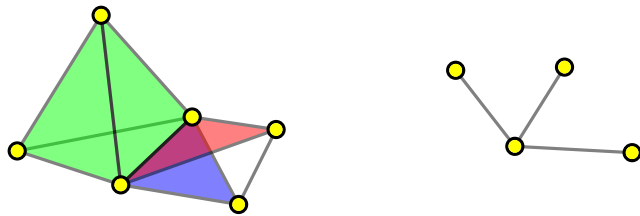
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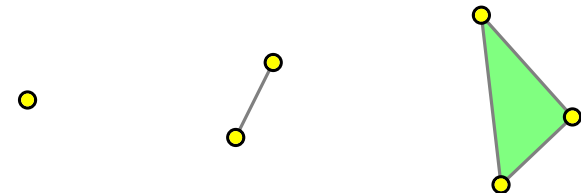


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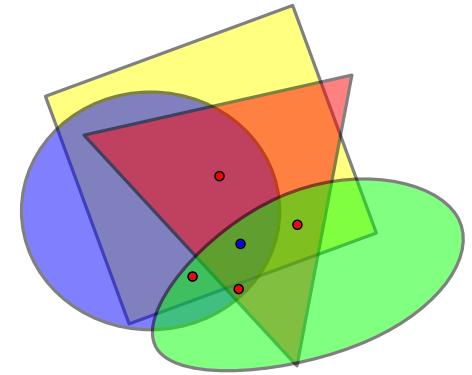
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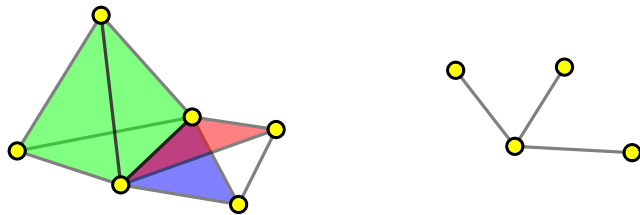
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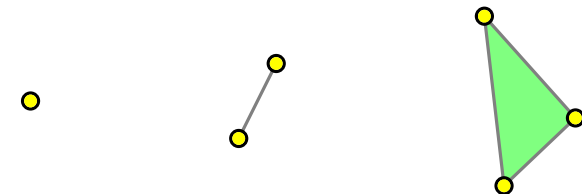


$\{\emptyset, \{1\}, \{2\}, \{3\}, \{4\}, \{1, 2\}, \{1, 3\}, \{1, 4\}\}$

map singletons to points
in general position in
 $\mathbb{R}^d$, d large enough

take convex hulls

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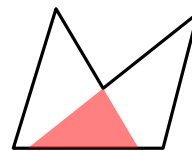


Why topological generalizations of Helly's theorem?

- ★ "Sort out" an industry of
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4 [Swanepoel 2003]



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Common eigen vector
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$\lfloor \frac{3d}{2} \rfloor$ [Polyanskii 2016]

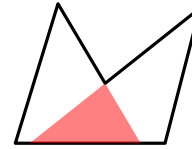
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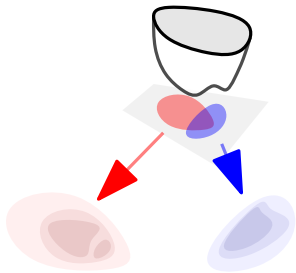
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Study the **combinatorial** aspects of optimization problems over **continuous** sets.



Given $f : \mathbb{R}^d \rightarrow \mathbb{R}$ and subsets $C_1, C_2, \dots, C_n \subset \mathbb{R}^d$, minimize f over $\bigcap_{i=1}^n C_i$.

Define $F_t = \{C_1, C_2, \dots, C_n\} \cup \{f^{-1}((-\infty, t))\}$.

$\inf_{\bigcap_{i=1}^n C_i} f = \inf \{t \in \mathbb{R} : \bigcap_{A \in F_t} A \neq \emptyset\}$.

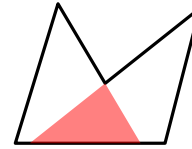
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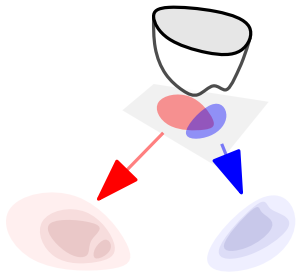
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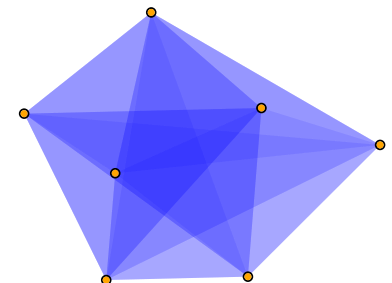
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- ★ Foundation of **combinatorial convexity**, for instance...

Fractional Helly Theorem. For any $0 < \alpha \leq 1$ and any finite family of n convex sets in $\mathbb{R}^d$, if $\alpha \binom{n}{d+1}$ of the $(d+1)$ -elements subsets of the family have non-empty intersection then some $(1 - (1 - \alpha)^{\frac{1}{d+1}})n$ elements of the family have non-empty intersection.

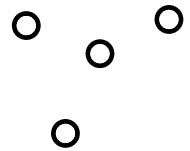
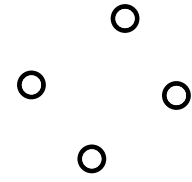
First selection lemma. For any finite point set $P \subset \mathbb{R}^2$ there exists a point $x \in \mathbb{R}^2$ contained in at least $\frac{2}{9} \binom{|P|}{3}$ triangles spanned by P .

Holds in any dimension with $\frac{2}{9}$ replaced by suitable constant.

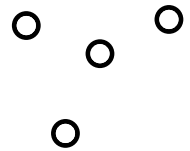
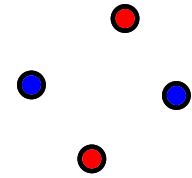


1. Non-embeddable complexes

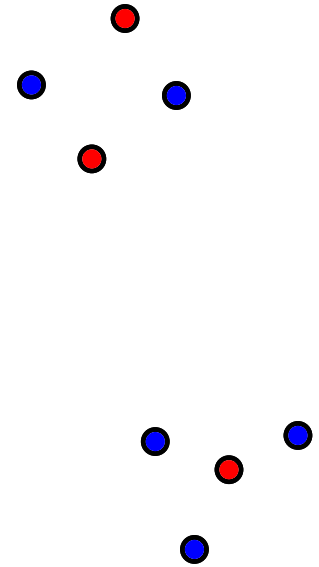
Radon's lemma. Any $n \geq d + 2$ points in $\mathbb{R}^d$ can be partitioned into two subsets whose convex hulls intersect.



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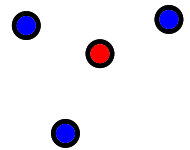
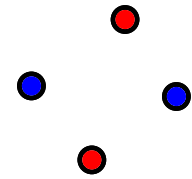
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Proof: Fix a frame, consider $x_1, x_2, \dots, x_n$ in $\mathbb{R}^d$ and define $\mathbf{x}_i := \begin{pmatrix} x_i \\ 1 \end{pmatrix} \in \mathbb{R}^{d+1}$.

Consider a non-trivial linear dependency $\sum_{i=1}^n \alpha_i \mathbf{x}_i = \vec{0}$.

Split it: $\sum_{i:\alpha_i \geq 0} \alpha_i \mathbf{x}_i = \sum_{i:\alpha_i < 0} -\alpha_i \mathbf{x}_i$.

Last coordinate ensures that $\sum_{i:\alpha_i \geq 0} \alpha_i = \sum_{i:\alpha_i < 0} -\alpha_i$. $\square$



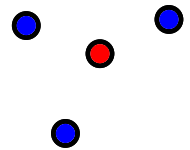
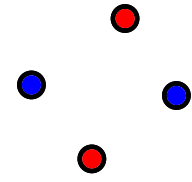
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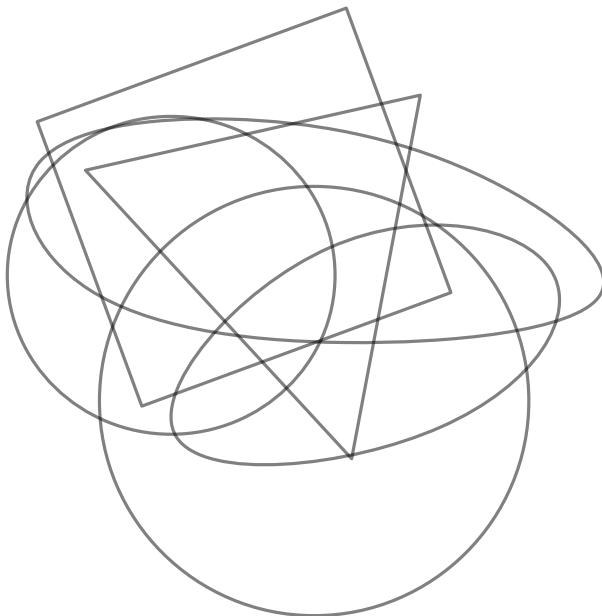
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Classical proof of Helly's theorem from Radon's lemma.

Let $\mathcal{F}$ be a family of convex sets in $\mathbb{R}^2$ where any three intersect.



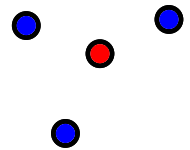
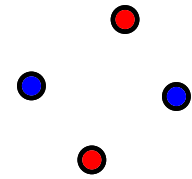
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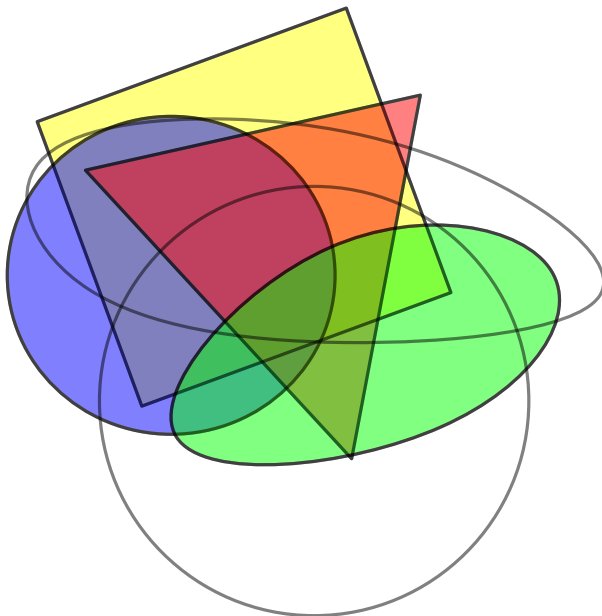
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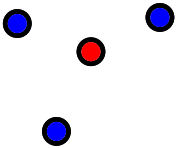
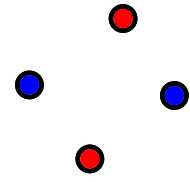
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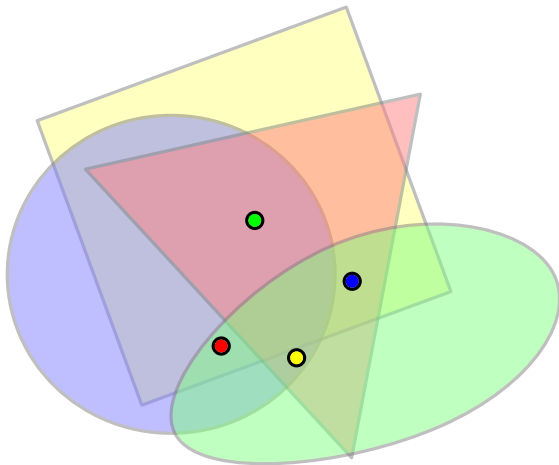


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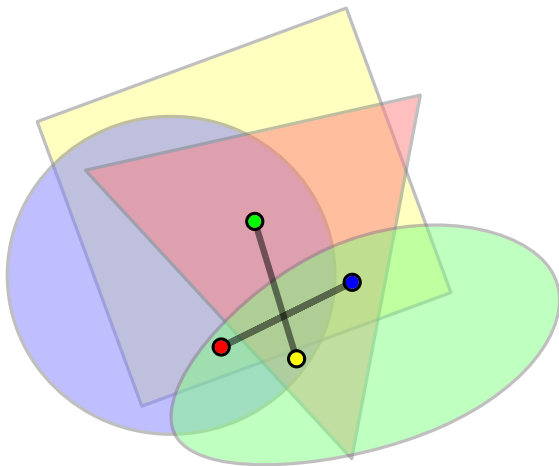
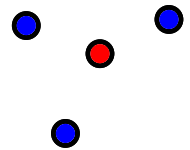
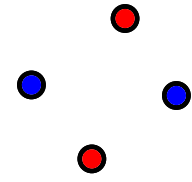
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Split it: $\sum_{i:\alpha_i \geq 0} \alpha_i \mathbf{x}_i = \sum_{i:\alpha_i < 0} -\alpha_i \mathbf{x}_i$.

Last coordinate ensures that $\sum_{i:\alpha_i \geq 0} \alpha_i = \sum_{i:\alpha_i < 0} -\alpha_i$. $\square$



Classical proof of Helly's theorem from Radon's lemma.

Let $\mathcal{F}$ be a family of convex sets in $\mathbb{R}^2$ where any three intersect.

Consider four of them $A_1, \dots, A_4$

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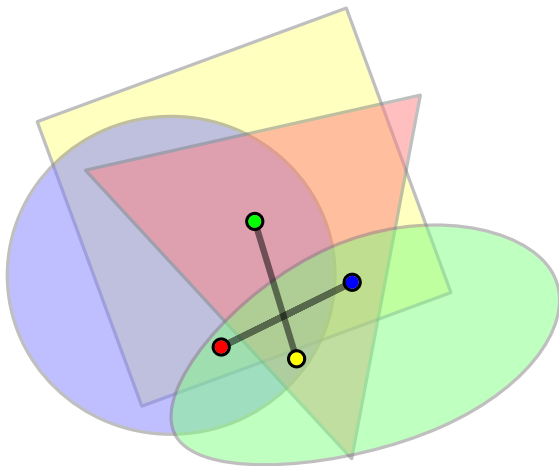
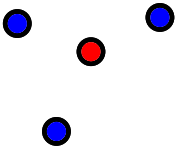
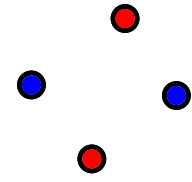
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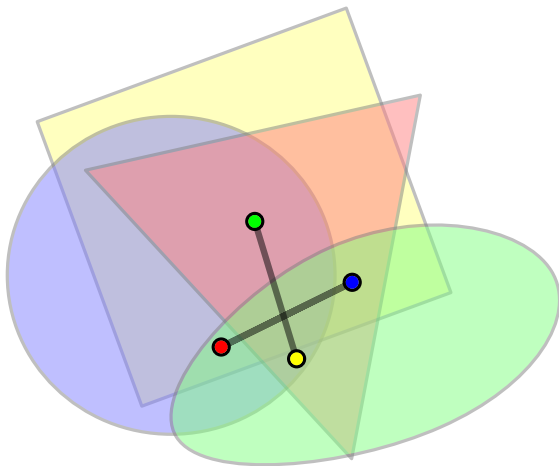
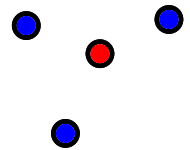
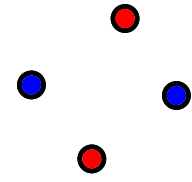
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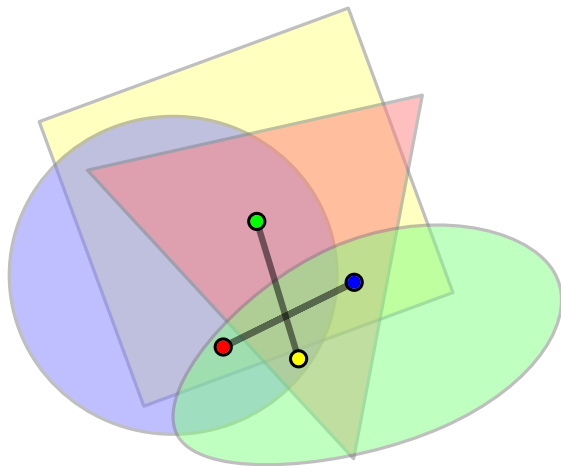
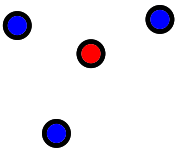
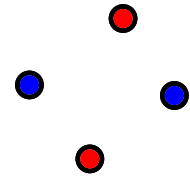
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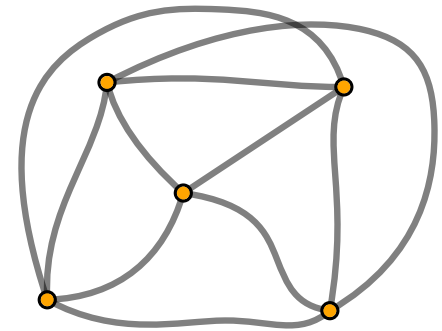
This is a **non-embeddability** argument.

Same argument, replacing Radon's lemma by the non-planarity of K_5 .

Fact. In any (topological) drawing of K_5 in the plane there are two vertex-disjoint edges whose images intersect.

K_5 is the complete graph on 5 vertices.

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edges to Jordan arcs joining the image of their vertices.*

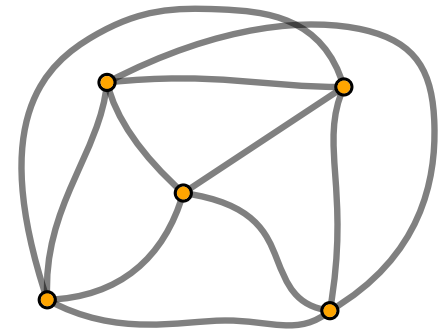


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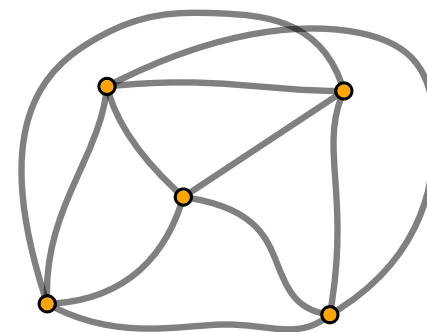
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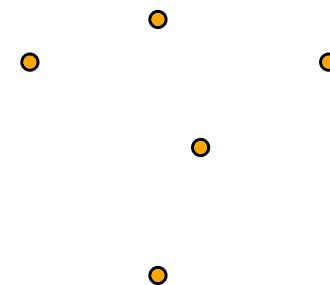
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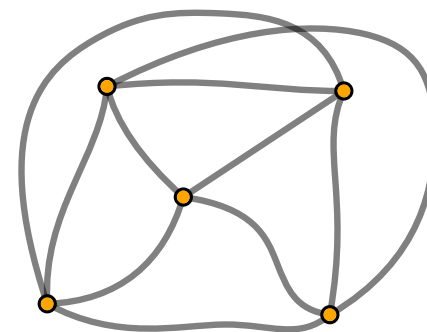


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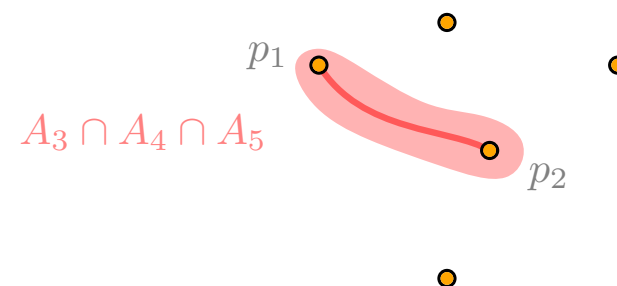
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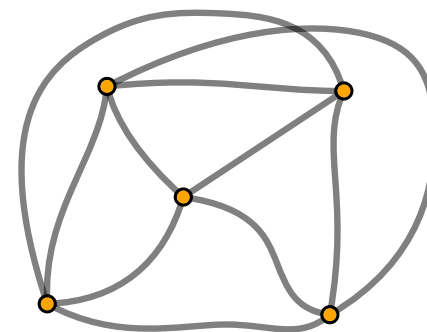


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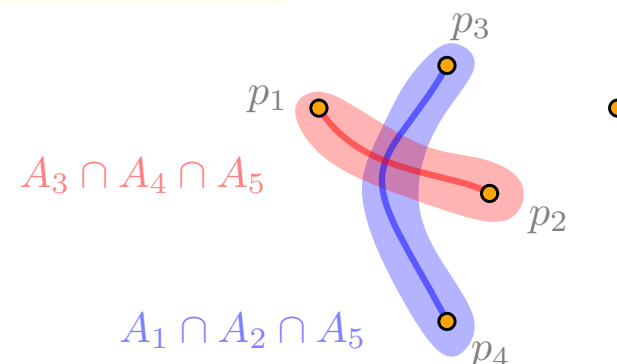


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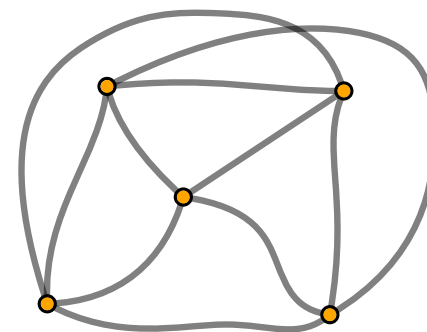


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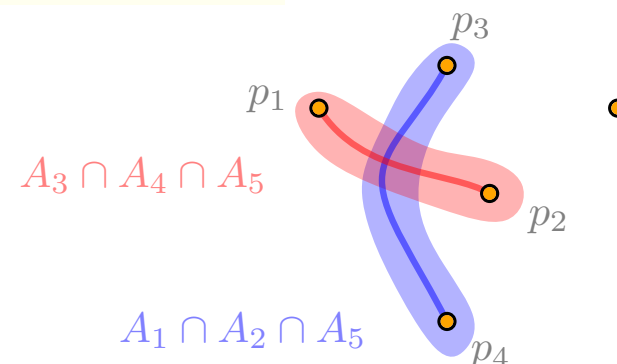
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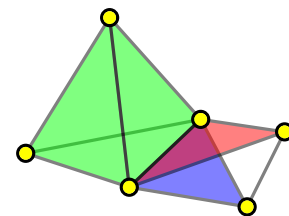


This idea generalizes to **higher dimension**.

Any abstract simplicial complex K has a **geometric realization** $|K|$.

Map singletons $x \in K$ to points $p(x)$ in general position in $\mathbb{R}^d$, d large enough.

*For every $\sigma \in K$ let $|\sigma| \stackrel{\text{def}}{=} \text{conv}\{p(x) : x \in \sigma\}$, and $|K| \stackrel{\text{def}}{=} \bigcup_{\sigma \in K} |\sigma|$.
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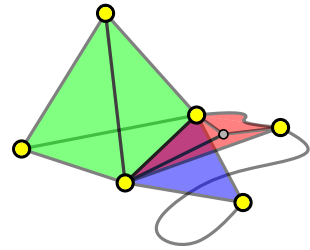
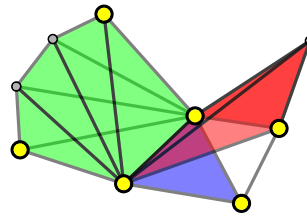
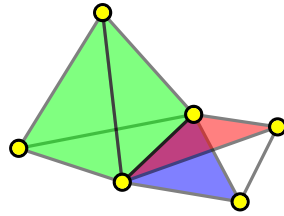
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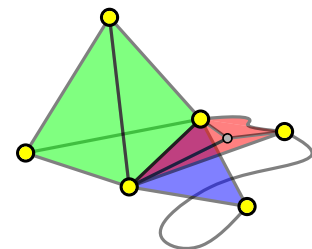
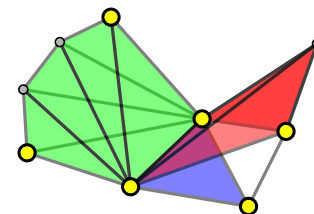
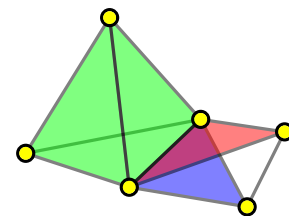
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[Van Kampen, 1932][Flores, 1933]

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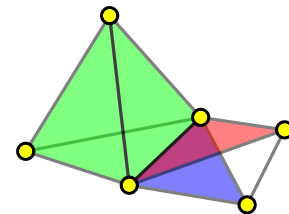


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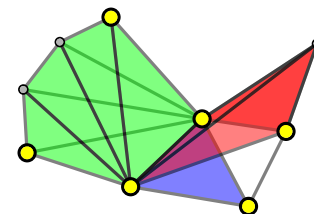
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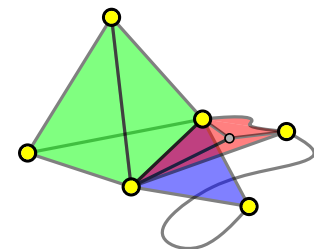
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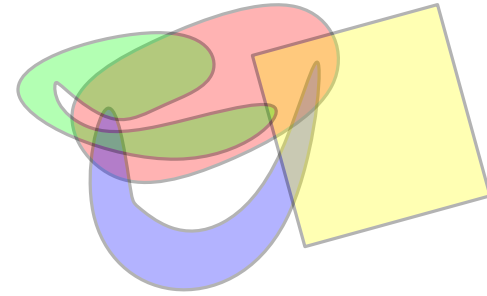
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This argument can be pushed a bit further (with help from Ramsey).

Now consider a family of subsets of $\mathbb{R}^2$ where every subfamily intersects in at most **two** path-connected components.



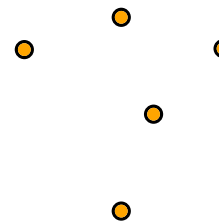
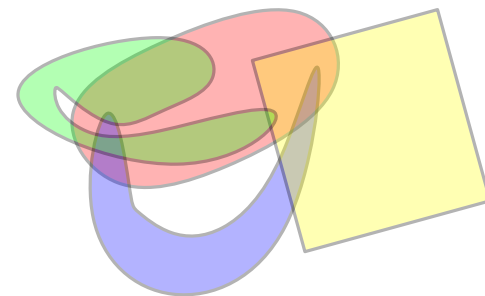
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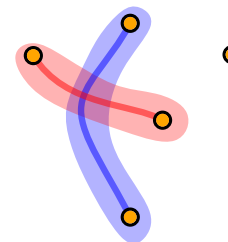
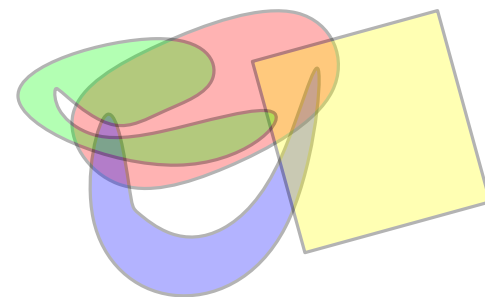
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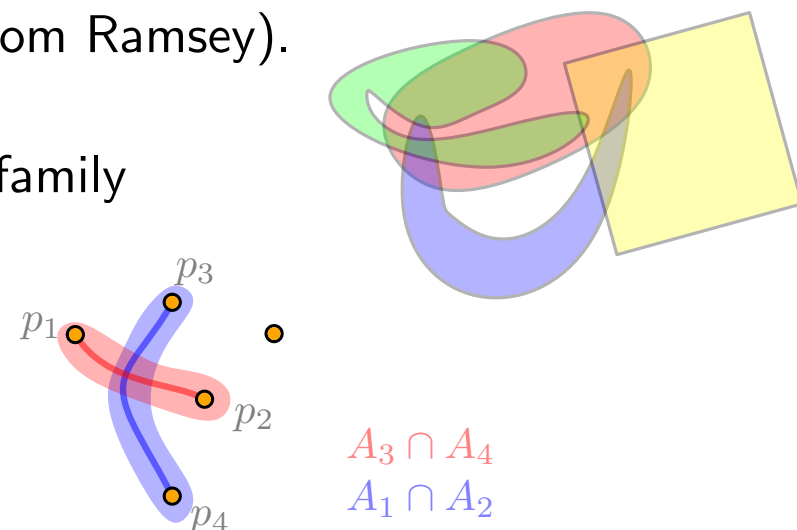
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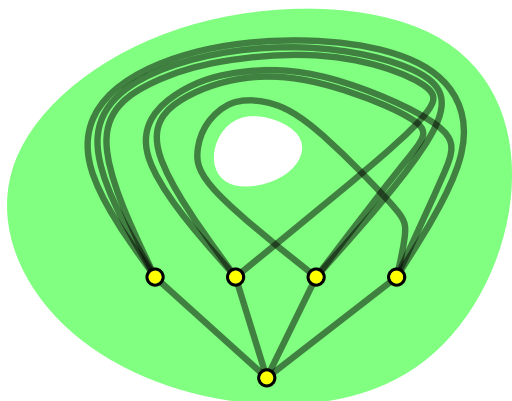
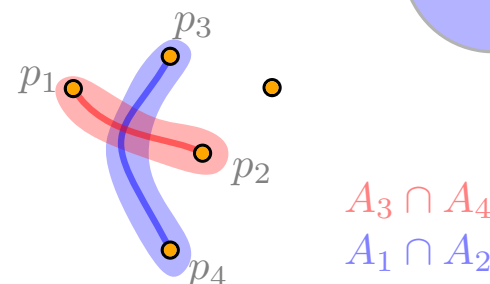
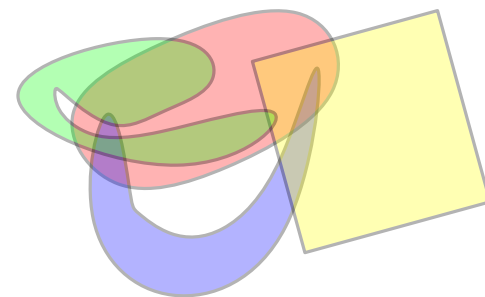
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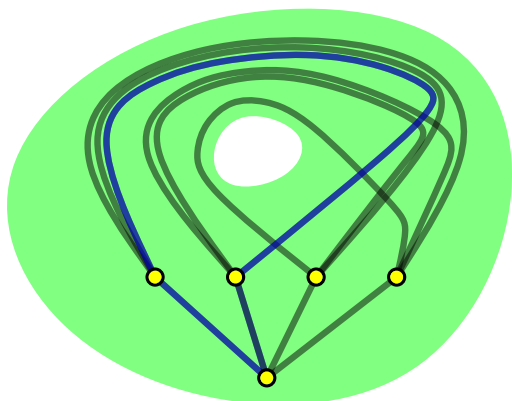
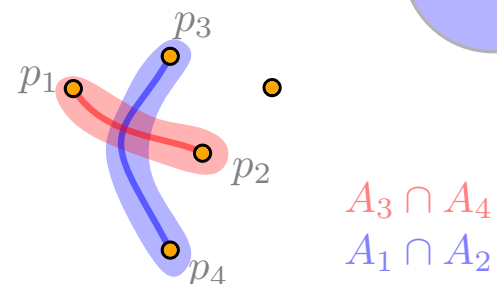
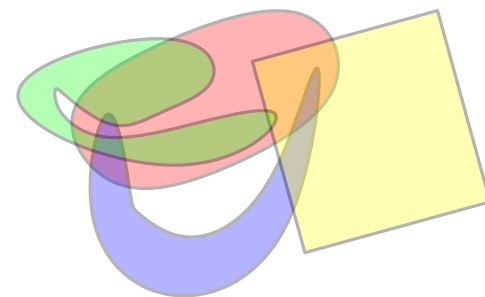
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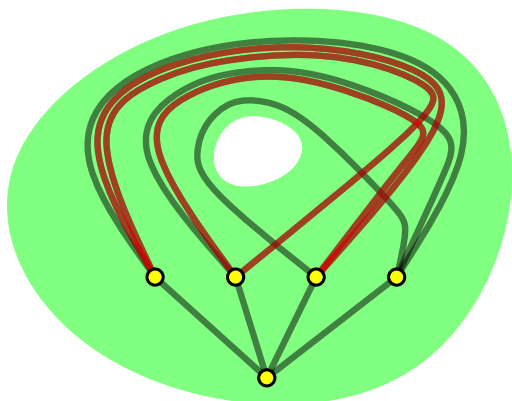
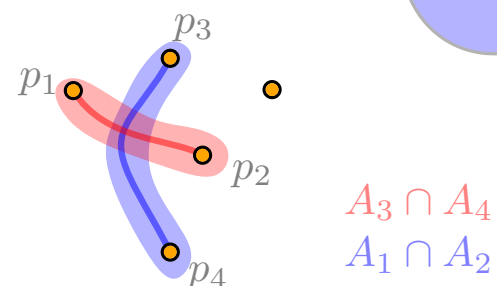
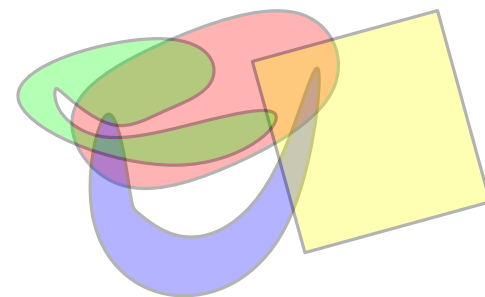
For any three indices u, v, w , some pair in p_u, p_v, p_w can be connected inside $\cap_{j \neq u, v, w} A_j$.

This draws a graph with no independent set of size 3.

So for k large enough, this graph contains a K_5 (Ramsey's theorem).

Not enough... but this type of idea works.

Ramsey for 3-uniform hypergraphs yields a K_5 where different edges miss different "dummy" subsets.



This argument can be pushed a bit further (with help from Ramsey).

Now consider a family of subsets of $\mathbb{R}^2$ where every subfamily intersects in at most **two** path-connected components.

Assume every k intersect (for some value of k).

So pick $k + 1$ of them: $A_1, A_2, \dots, A_{k+1}$.

Pick $p_j \in \cap_{i \neq j} A_i$

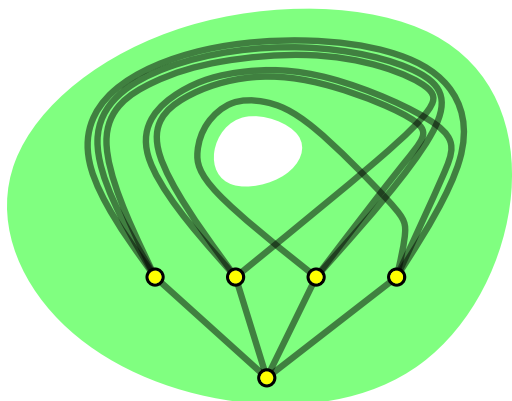
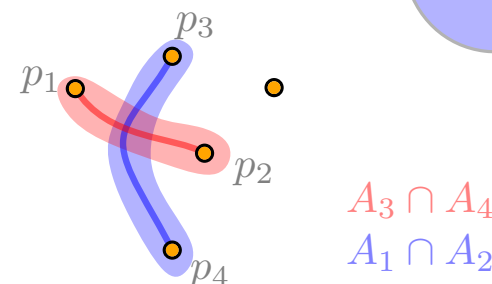
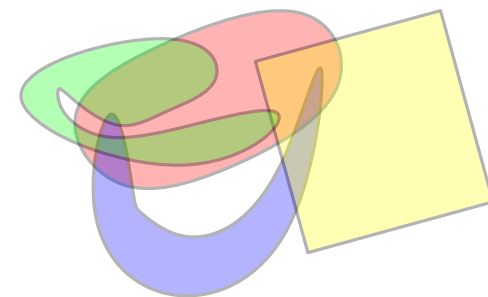
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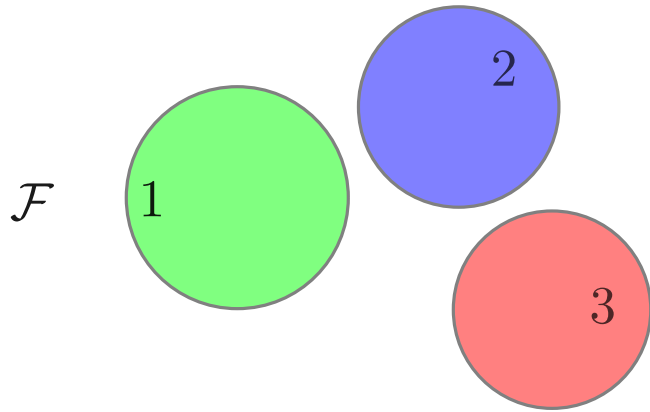
And quite a bit further (using homological relaxations of embeddings + Ramsey on homology classes).

Theorem. The Helly number of any family of subsets of $\mathbb{R}^d$ can be bounded by a function of d and

$$\max_{G \subseteq \mathcal{F}} \max_{i \leq \lceil \frac{d}{2} \rceil - 1} \tilde{\beta}_i(\cap G, \mathbb{Z}_2).$$

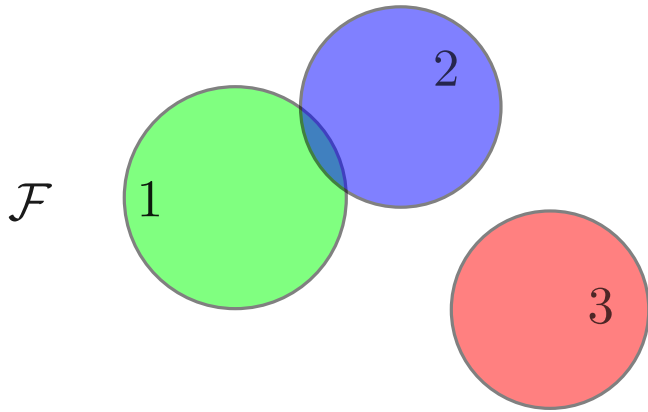
2. Nerve complexes

The **nerve** $\mathcal{N}(\mathcal{F})$ of a family $\mathcal{F}$ of sets is $\mathcal{N}(\mathcal{F}) = \{\mathcal{G} : \mathcal{G} \subseteq \mathcal{F} \text{ and } \cap \mathcal{G} \neq \emptyset\}$.



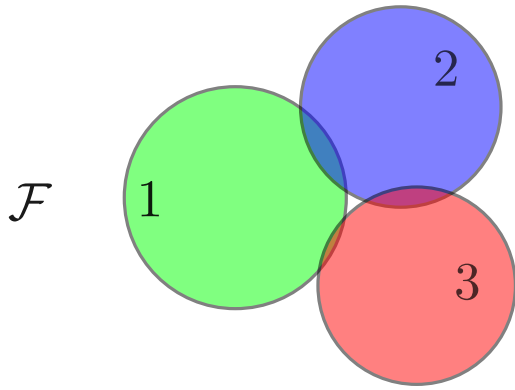
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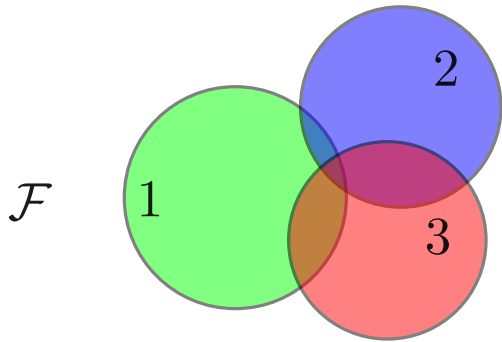
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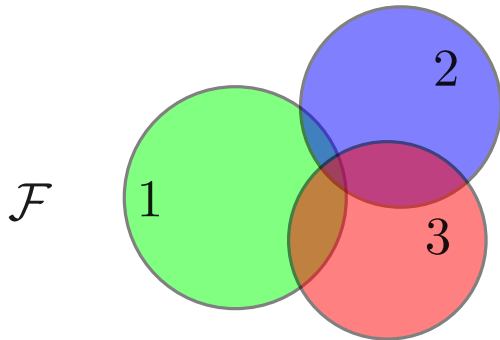
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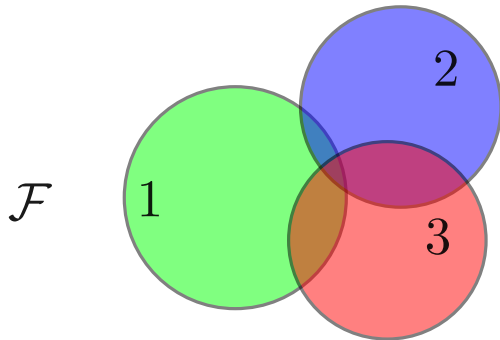
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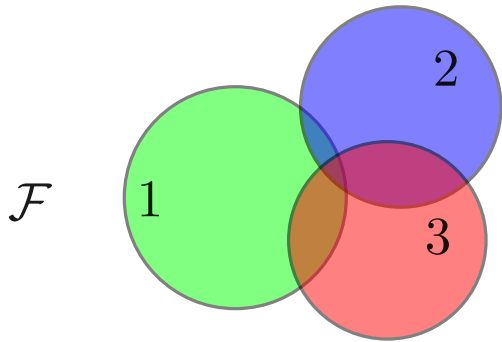
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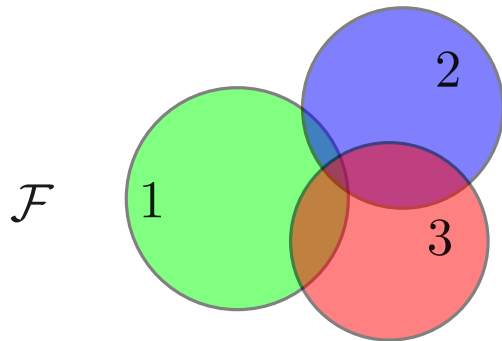
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so...

Families eluding Helly-type theorems have nerves with "simplicial holes" of arbitrary high dimensions.

Statement about collections of families, eg. "all finite families of convex sets in $\mathbb{R}^2$ ".

Dimension of simplicial holes VS size of family.

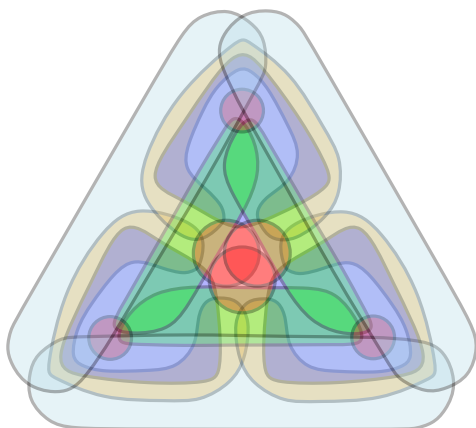
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non-collapsible subcomplexes (discrete homotopy)...*

One tool available for **nerves** of "nice" families.



Families of sets where subfamilies have "hole-free" intersections.

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Contractible = homotopic to a point.

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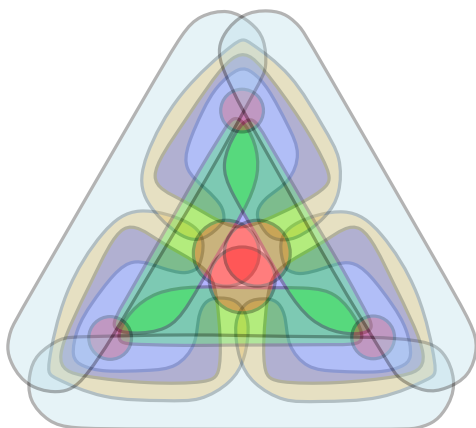
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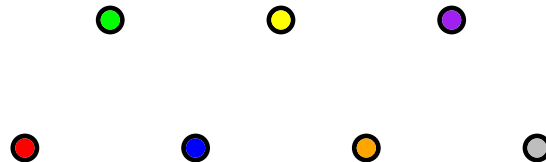
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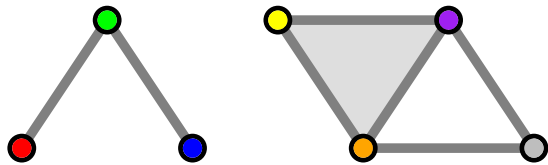
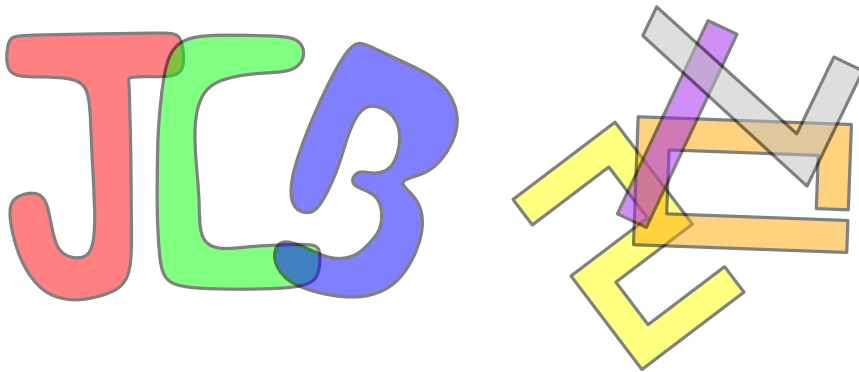
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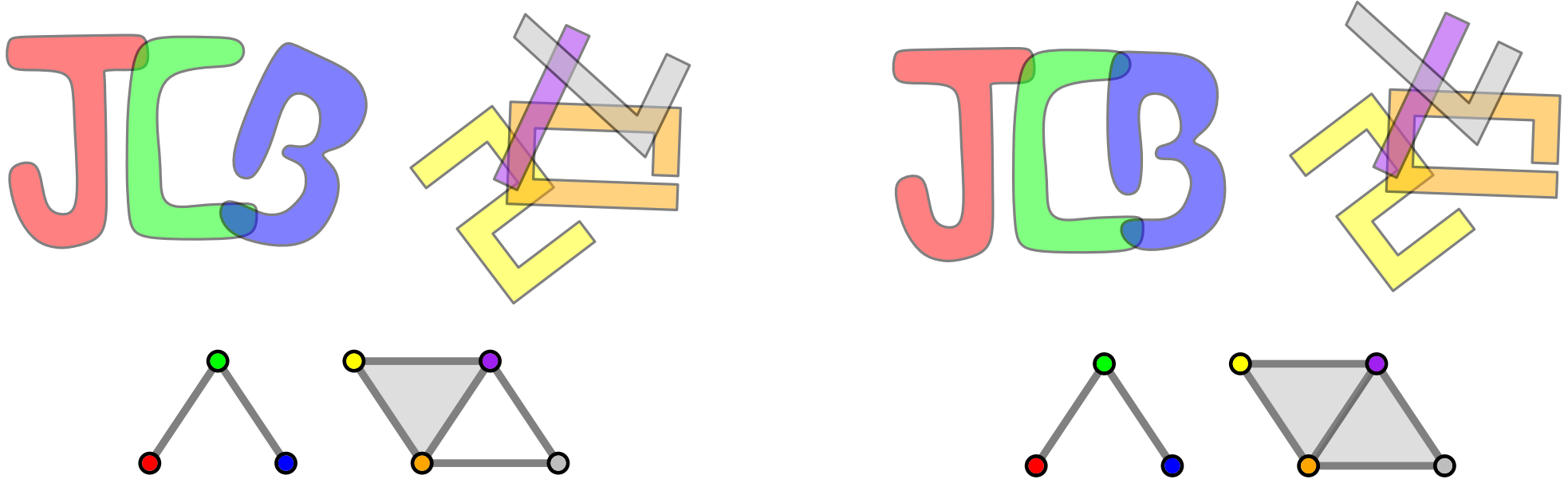
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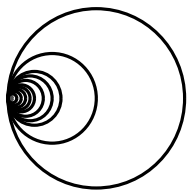
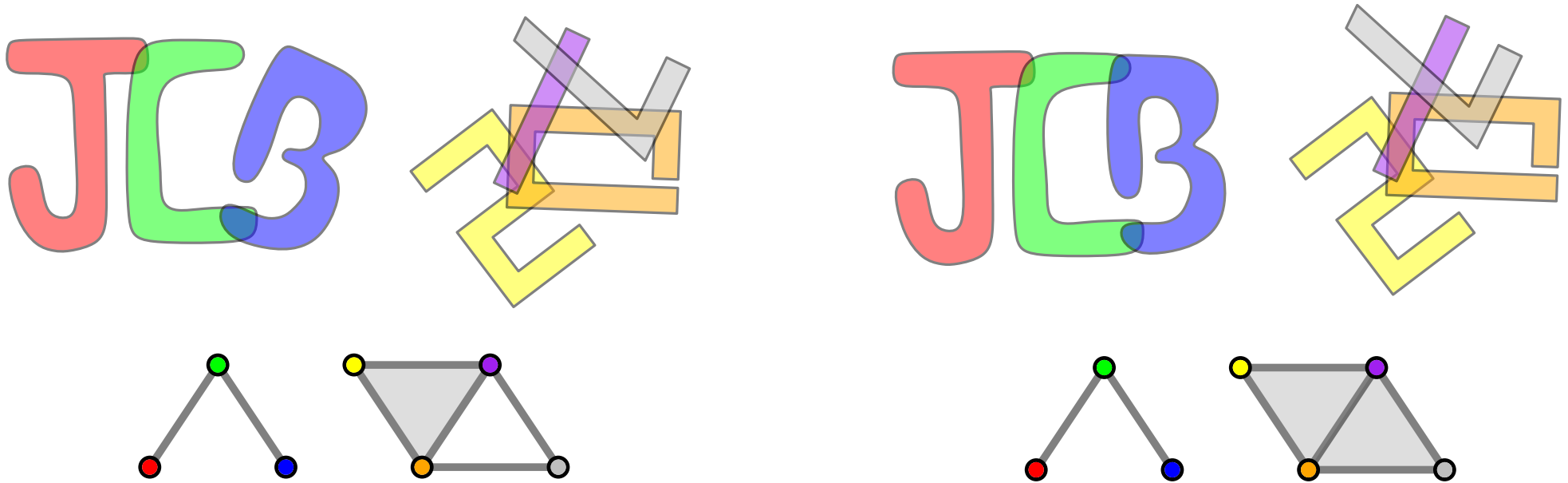
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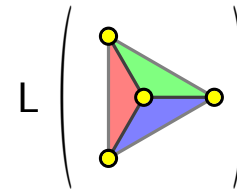
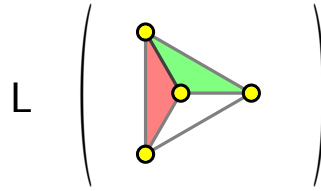
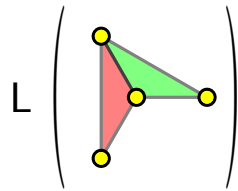
Fact. Any open subset $\mathbb{R}^d$ has trivial $\mathbb{Q}$ -homology in dimension $> d$.

Homological Helly Theorem. Let $\mathcal{F}$ be an acyclic cover in $\mathbb{R}^d$. If every $d + 1$ elements of $\mathcal{F}$ have nonempty intersection then $\mathcal{F}$ has nonempty intersection.

The **Leray number** of a simplicial complex K with vertex set V is

$$L(K) = \min\{\ell \in \mathbb{N} : \forall i \geq \ell, \forall S \subseteq V, \quad \tilde{H}_i(K[S], \mathbb{Q}) = 0\}$$

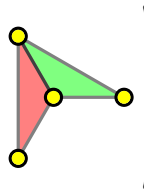
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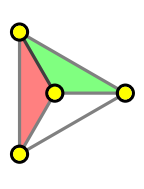


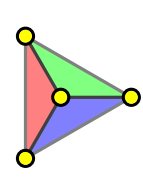
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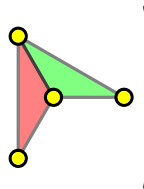
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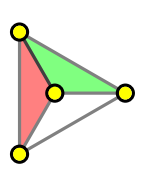
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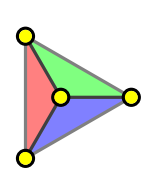
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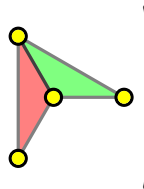
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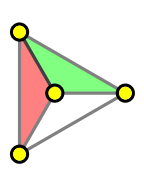
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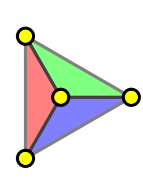
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Nerve theorem $\Rightarrow$ Nerves of acyclic covers in $\mathbb{R}^d$ have Leray number $\leq d$.

Induced subcomplex of a nerve = nerve of a subfamily.

Theorem. If $L(K) \leq d$ then for any $r \geq 0$,

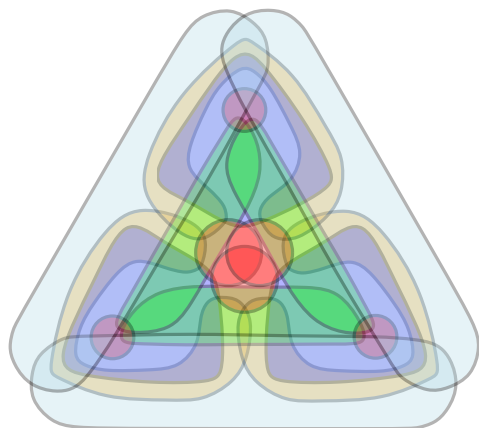
$$f_d(K) > \binom{n}{d+1} - \binom{n-r}{d+1} \Rightarrow f_{d+r}(K) > 0.$$

[Eckhoff 1984] [Kalai 1984]

$n = \#$ vertices and $f_i(K) = \#$ i -dimensional faces of K .

Many low-dimensional faces $\Rightarrow$ existence of high-dimensional faces.

Sharp: r copies of $\mathbb{R}^d$ and $n - r$ hyperplanes in general position.



Excerpt from the theory of combinatorial convexity.

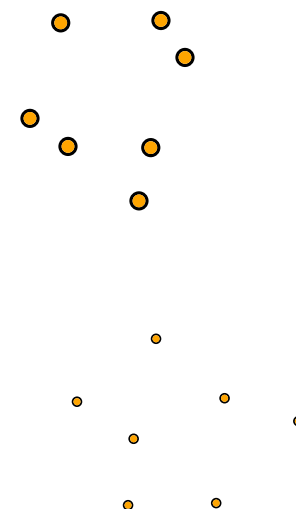
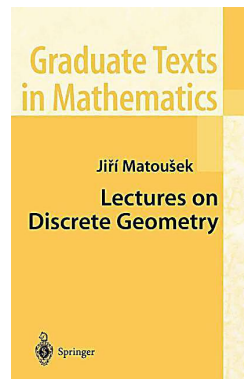
Fractional Helly theorem. For any $0 < \alpha \leq 1$ there exists $\beta > 0$ such that any finite family $\mathcal{F}$ of convex sets of $\mathbb{R}^d$ some $\alpha \binom{\mathcal{F}}{d+1}$ of the $(d+1)$ -tuples intersect, then some $\beta|\mathcal{F}|$ of the sets intersect.

Tverberg's theorem. Any $(r-1)(d+1)+1$ points in $\mathbb{R}^d$ can be partitioned into r parts so that the convex hulls of these parts intersect.

First selection lemma. For any $d \geq 1$ There exists $c_d > 0$ such that for any set P of n points in $\mathbb{R}^d$, a fraction c_d of the d -simplices spanned by P intersect.

Weak ϵ -nets theorem. For every d and $\epsilon > 0$ there exists a constant $f(d, \epsilon)$ such that for any finite point set P in $\mathbb{R}^d$, there exists a subset $N \subseteq \mathbb{R}^d$ with $|N| \leq f(d, \epsilon)$ such that any convex set C with $|C \cap P| \geq \epsilon|P|$ must intersect N .

(p, q) theorem. For any family of convex sets in $\mathbb{R}^d$ where among any p , some q intersect, there exist $c(p, q, d)$ points that intersect every member of the family.



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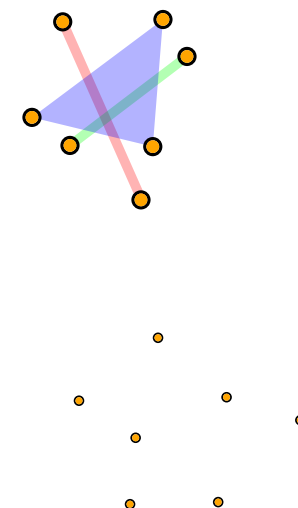
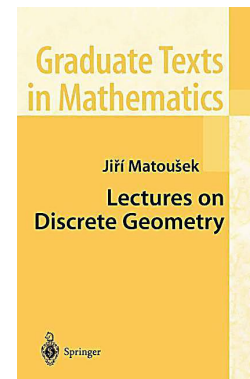
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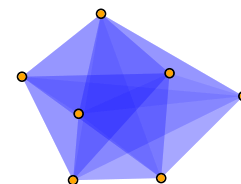
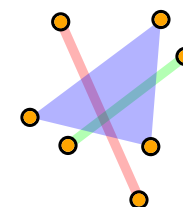
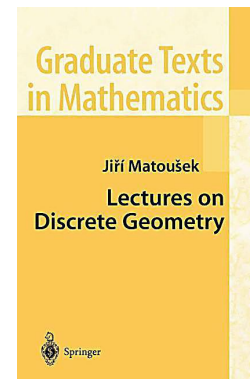
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Weak ϵ -nets theorem. For every d and $\epsilon > 0$ there exists a constant $f(d, \epsilon)$ such that for any finite point set P in $\mathbb{R}^d$, there exists a subset $N \subseteq \mathbb{R}^d$ with $|N| \leq f(d, \epsilon)$ such that any convex set C with $|C \cap P| \geq \epsilon|P|$ must intersect N .

(p, q) theorem. For any family of convex sets in $\mathbb{R}^d$ where among any p , some q intersect, there exist $c(p, q, d)$ points that intersect every member of the family.



Excerpt from the theory of combinatorial convexity.

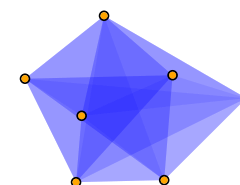
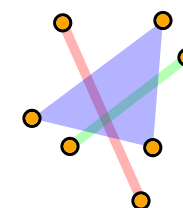
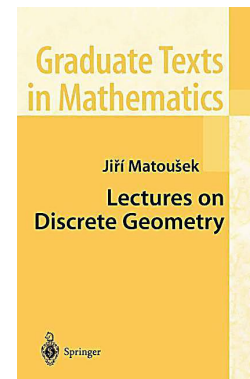
Fractional Helly theorem. For any $0 < \alpha \leq 1$ there exists $\beta > 0$ such that any finite family $\mathcal{F}$ of convex sets of $\mathbb{R}^d$ some $\alpha \binom{\mathcal{F}}{d+1}$ of the $(d+1)$ -tuples intersect, then some $\beta|\mathcal{F}|$ of the sets intersect.

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The Kalai-Eckhoff upper bound implies the fractional Helly theorem.

With constant $\beta = 1 - (1 - \alpha)^{\frac{1}{d+1}}$, sharp exponent, only known proof.

All five theorems generalize to families of sets with d -Leray nerves.

Replace "convex hull of X " by "intersection of all family members containing X ".

Theorem. Let $\mathcal{F}$ be a family of subsets of $\mathbb{R}^d$ such that every subfamily intersects in at most r connected components, each of which is acyclic. $L(N(\mathcal{F})) \leq rd - 1$.

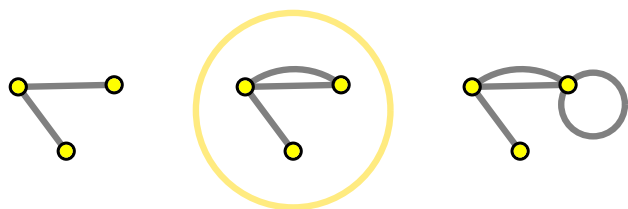
Helly-type theorem	}	for acyclic families with bounded number of connected components of intersection
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Fix the nerve by adding multiple simplices in case of multiple components

Simplicial poset: every lower-interval is isomorphic to a the face lattice of a simplex.



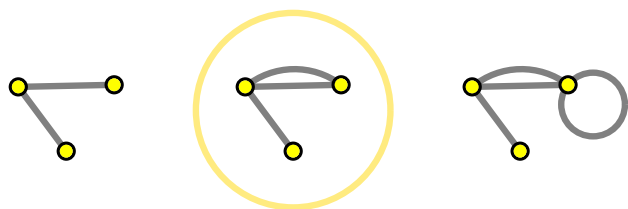
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Arguments about co-faces may not extend

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Extend nerve theorem.

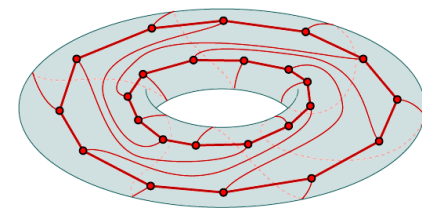
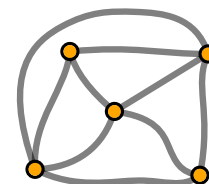
Analyze the removal of multiplicities $\simeq$ projection with fibers of bounded size.

3. Some perspectives

Graphs are often used to model pairwise interactions, and hypergraphs are natural generalization to model many-wise interactions.

There is a rich theory of **geometric graphs**.

$$V - E + F = 2, E \leq 3V - 6, Cr(G) \geq \frac{E^3}{64V^2} \dots$$



The natural generalization of geometric graphs are **simplicial complexes**.

Non-planarity of $K_5 \rightsquigarrow$ Van Kampen-Flores theorem.

Conjecture. Let K be a d -dimensional simplicial complex on n vertices. If there exists an injection $|K| \rightarrow \mathbb{R}^{2d}$ then K has $O(n^d)$ simplices.

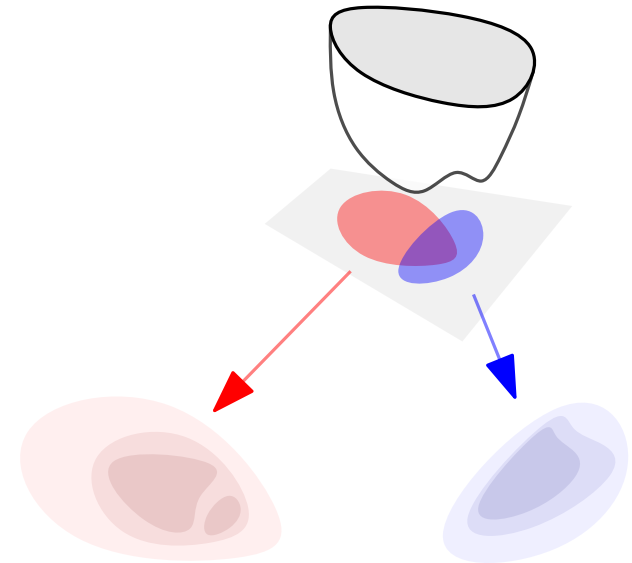
[Grünbaum][Kalai] [Sarkaria] [Dey]

Open for any $d \geq 2$.

Coming back to one of the original motivations...

Theorem. The Helly number of any family of subsets of $\mathbb{R}^d$ can be bounded by a function of d and

$$\max_{G \subseteq \mathcal{F}} \max_{i \leq \lceil \frac{d}{2} \rceil - 1} \tilde{\beta}_i(\cap G, \mathbb{Z}_2).$$



$\Rightarrow$ Optimization problems with **bounded homological complexity** are "solvable" in linear time.

In terms of black boxes such as "minimize f over $O(1)$ of the constraints".

Level sets of constraints have bounded homology in dimension $< \lceil \frac{d}{2} \rceil$.

Thank you for your attention